Check-In

Review: Machine-Code Optimization

Give an example of a snippet of x64 code that benefits from two peephole optimizations

```
~~jmp LBL_1~~
LBL_1: ~~pushq %rax~~      nop
~~popq %rax~~
```

Admin

Announcements and Housekeeping

All Quizzes Graded!

Announcements and Housekeeping

hooray

Quiz 2

Announcements and Housekeeping

P	6%
A	24%
B	28%
C	15%
D	9%
F	19%

Quiz 2

Announcements and Housekeeping

P	27%
A	30%
B	15%
C	9%
D	6%
F	13%

Drew Davidson | University of Kansas

EECS 665

COMPILER CONSTRUCTED

[y] = randBool()

Flowgraphs

Previously…

Machine Code Optimization

Machine code optimization overview

Improving data allocation

- Register allocation

Improving Final Code

- Peephole optimization
- Instruction Pipelines

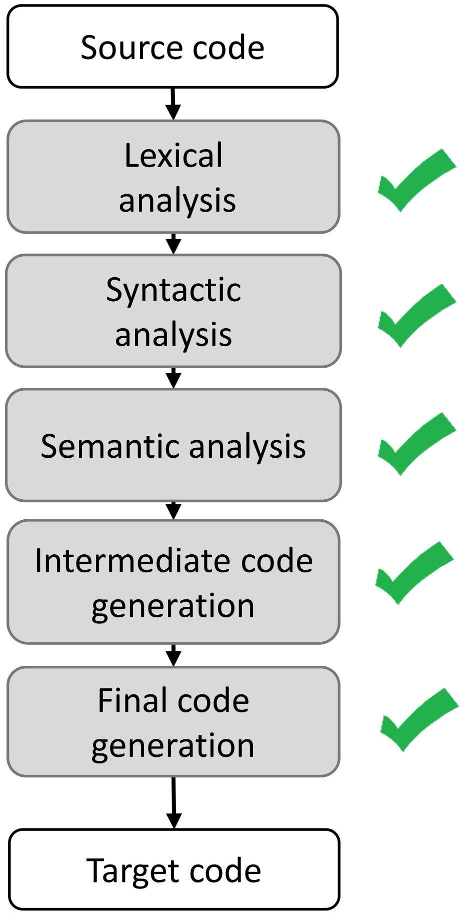
<u>You should know</u>

- Interference graphs
- Sharing AR slots / registers for allocation
- How/where to apply peephole optimizations
- How/where instruction reordering might aid a simple instruction pipeline

Optimization

Compiler Construction

Progress Pics



Basic source to target workflow:

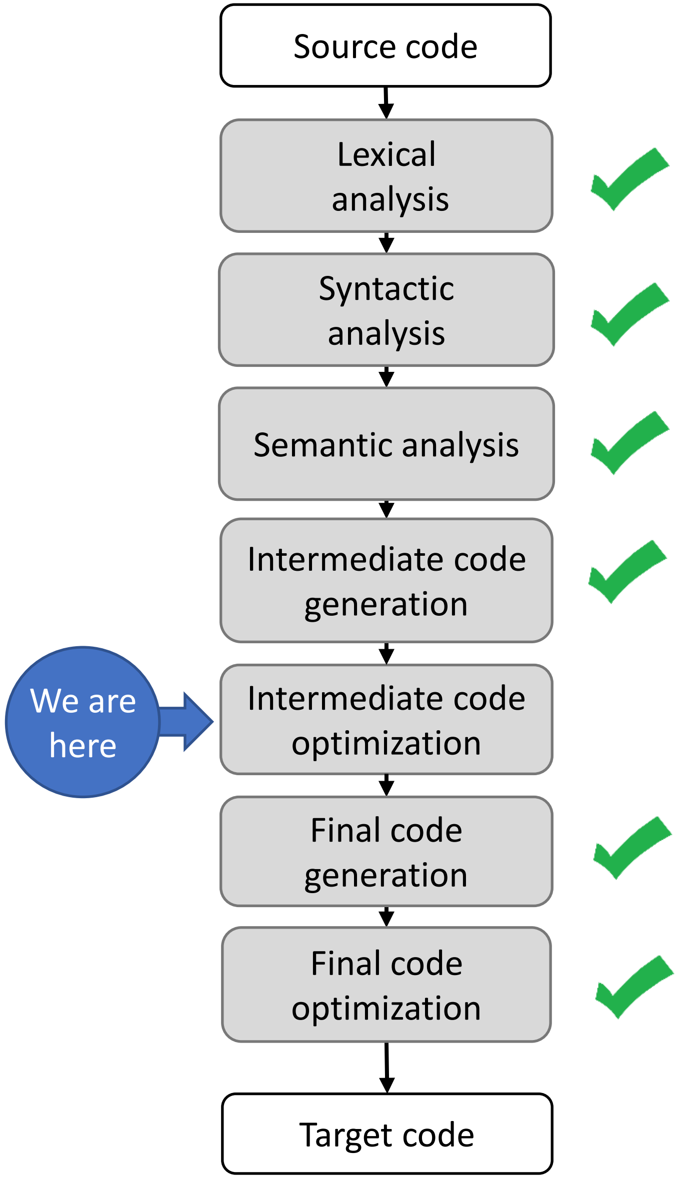
- Complete
- Outputs naïve code

Advanced workflow:

- “Postprocess” the output of a naïve phase

Compiler Construction

Progress Pics



Basic source to target workflow:

- Complete
- Outputs naïve code

Advanced workflow:

- "Postprocess" the output of a naïve phase
  - Discussed: final code "cleanup"
  - Next up: intermediate code

Lecture Outline

Flowgraphs

Program analysis:

- Goals
- Control flow graphs

Local Optimizations

- Dead code elimination
- Common subexpression elimination
- Constant/copy propagation

Optimization

Making faster IR programs

Flowgraphs: Program analysis

General constraints:

- We can't violate program semantics
- Minimal architecture details

Constraint-friendly goals:

- Don't do *useless* computation
- Don't do *redundant* computation

Simple Example: Constant Folding

Statically compute known expressions

- Replace the runtime expression with its value

Before
[z] := 1 + 2

Analysis
- Identify constant expressions
- Compute known value

Rewrite
- Replace expression with value

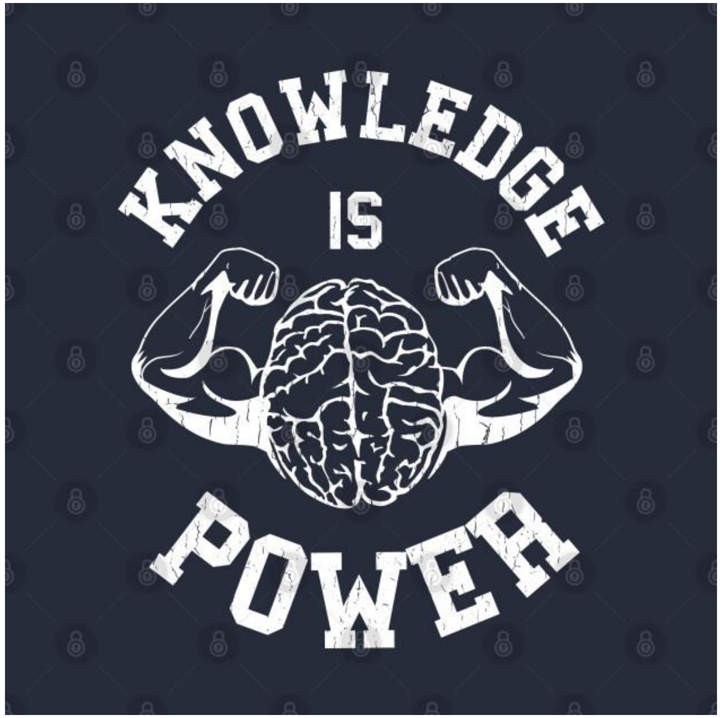
After
[z] := 3

Program Analysis

Flowgraphs: Program analysis

The more we know about the program the more we can improve it

- What might we be interested in knowing…?



“Structural” Properties of a Program

E.g. for a given program point:

- What paths lead there?
- Is it in a deeply nested loop?
- Is it reachable at all?

Knowing the above info supports other analyses

- Might a variable be uninitialized?

"Structural" Properties of a Program

E.g. for a given program point:

- What paths lead there?
- Is it in a deeply nested loop?
- Is it reachable at all?

We need a program abstraction to capture these details

Knowing the above info supports other analyses

- Might a variable be uninitialized?

Intuition: Flow charts

Flowgraphs: Program analysis

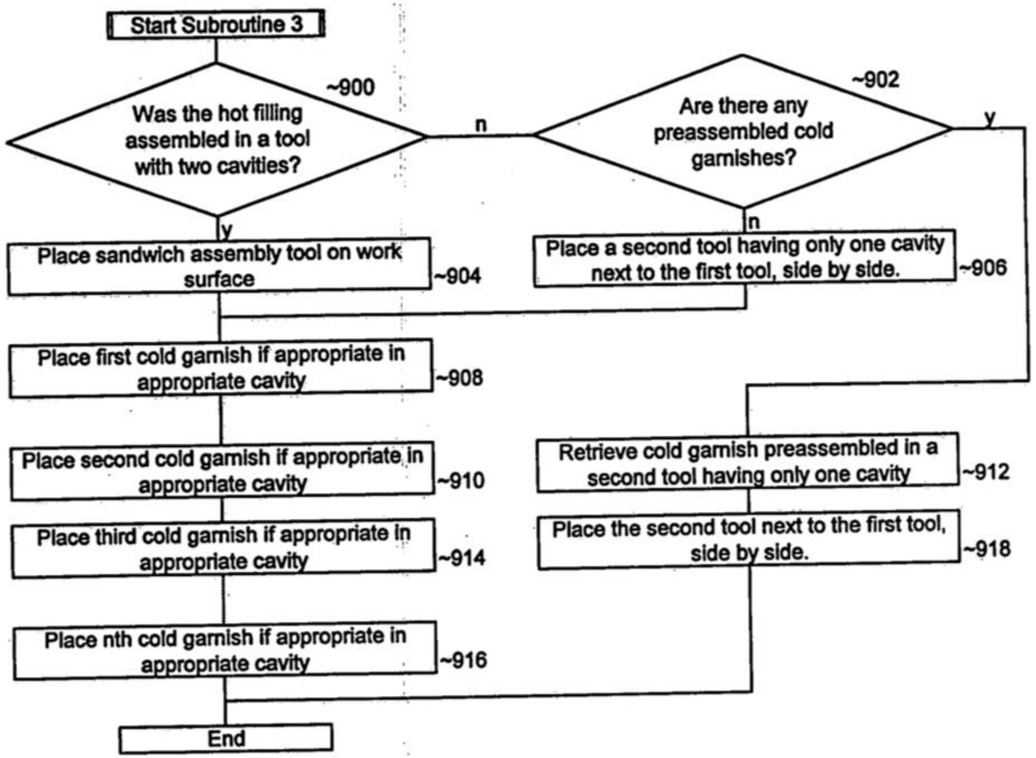
Notation

- Nodes are instructions
- Edges go to successor nodes

Operation

- Execute current instruction
- Proceed to the right successor

Fig. 59



Flow chart for building a sandwich, appearing in a McDonald's patent

Intuition: Flow charts

Flowgraphs: Program analysis

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- Nodes are instructions
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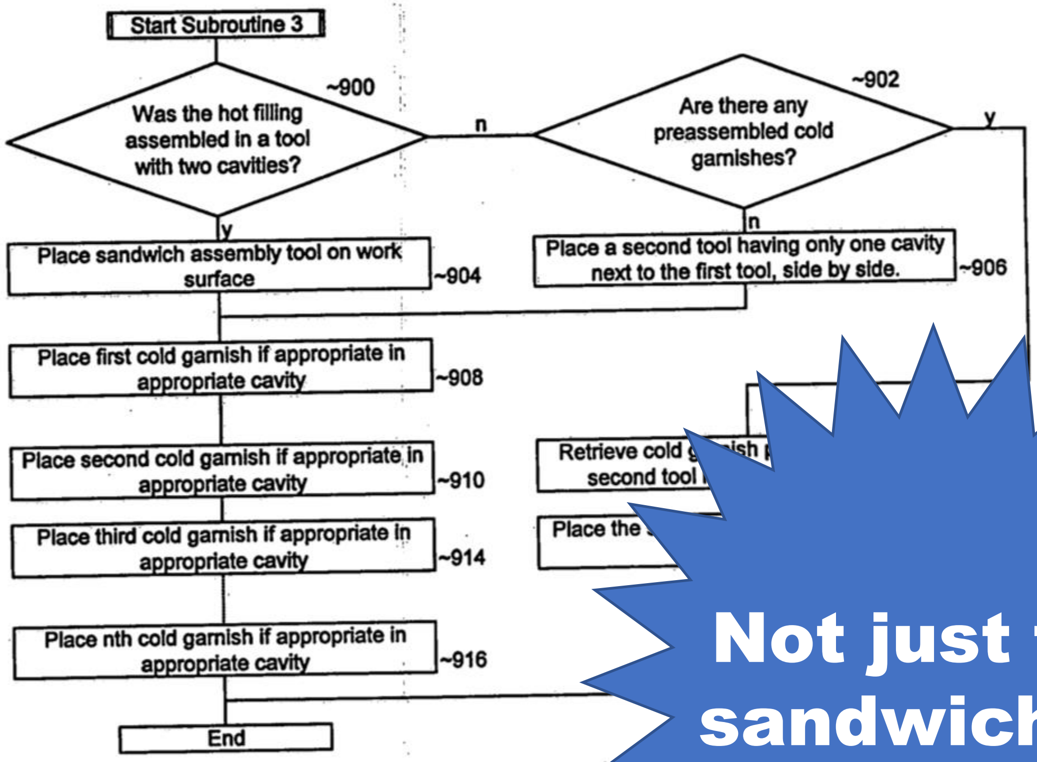
Fig. 59



Not just for sandwiches!

Flow chart for building a s... appearing in a McDonald's pa...

Intuition: Flow Charts … for Code?!

Flowgraphs: Program analysis

Notation

- Nodes are instructions
- Edges go to successor nodes

Operation

- Execute current instruction
- Proceed to the right successor

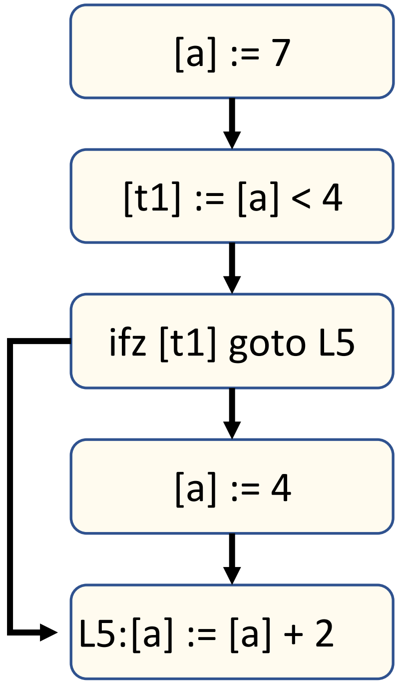
src code

```
a = 7;
if (a < 4){
  a = 4;
}
a += 2;
```

3AC code

```
    1. [a] := 7
    2. [t1] := [a] < 4
    3. ifz [t1] goto L5
    4. [a] := 4
L5: 5. [a] := [a] + 2
```

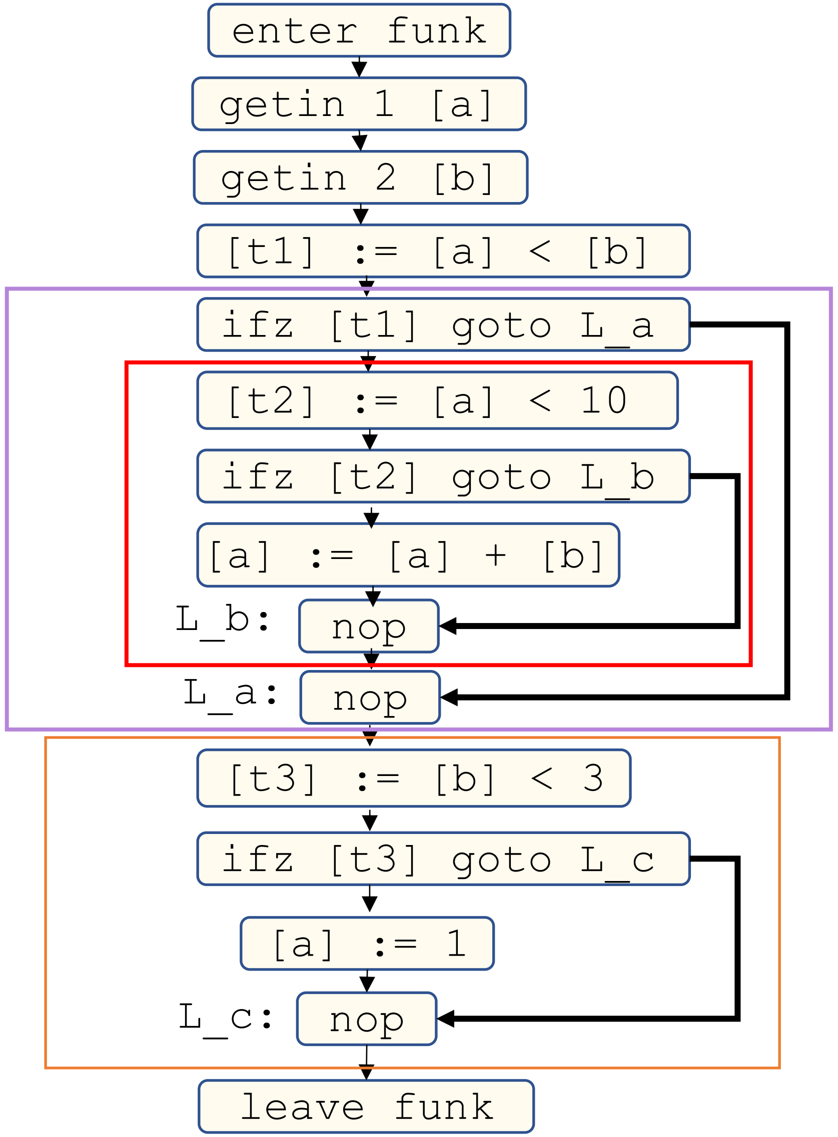
Instruction Flowgraph



Intuition: Flow Charts … FROM Code?!

Flowgraphs: Program analysis

```
void funk(int a, int b){
    if (a < b){
        if (a < 10){
            a = a + b;
        }
    }
    if (b < 3){
        a = 1;
    }
}
```



Intuition: Flow Charts … FROM Code?!

Flowgraphs: Program analysis

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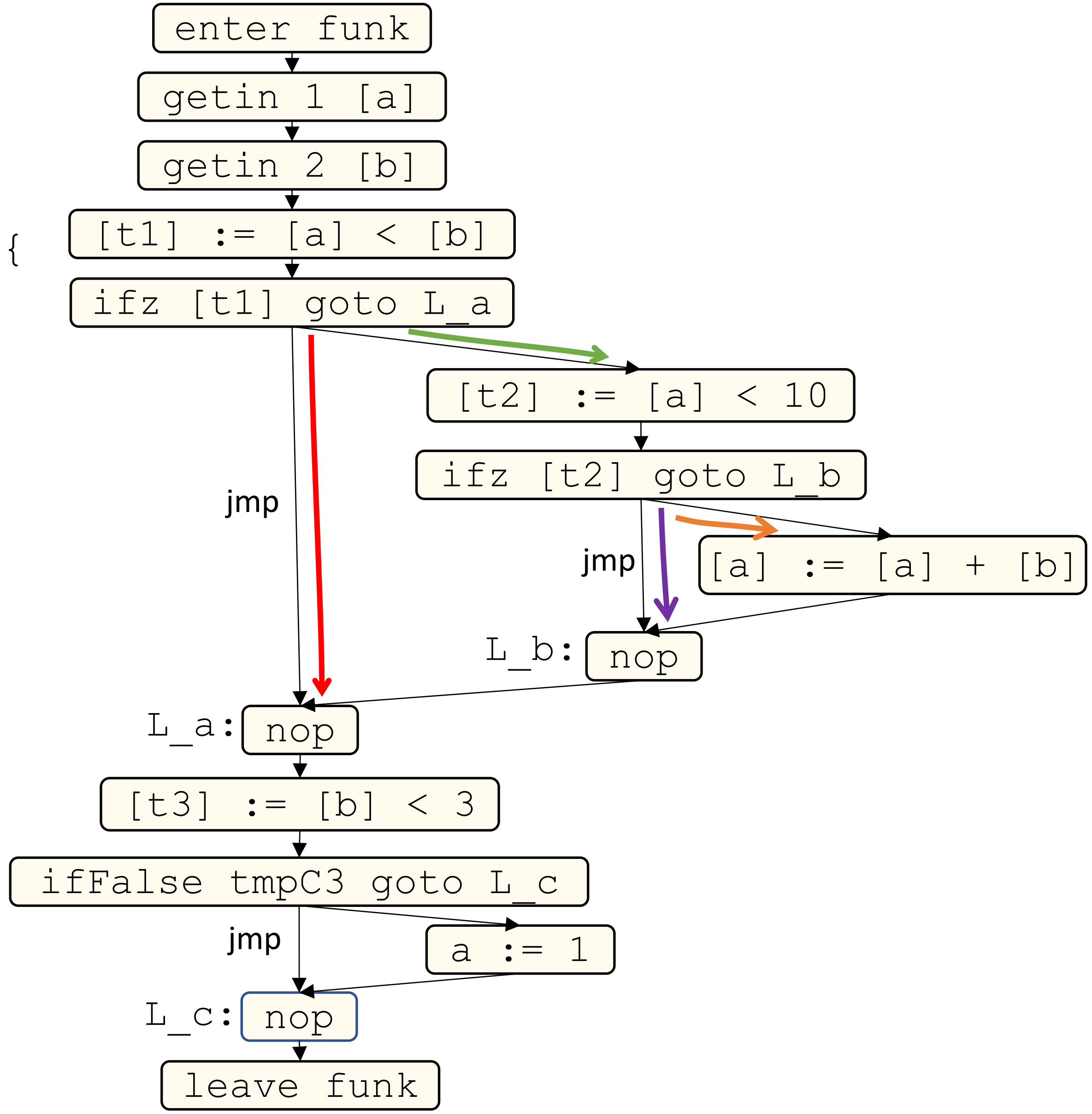


Code Flowcharts: Seem Familiar?

Flowgraphs: Program analysis

Maybe this is how you learned to think about code!

- It's a nice way to visualize the *control flow* of the program
- We can extend this intuition for program analysis

Lecture Outline

Flowgraphs

Program analysis:

- Goals
- Control flow graphs

Local Optimizations

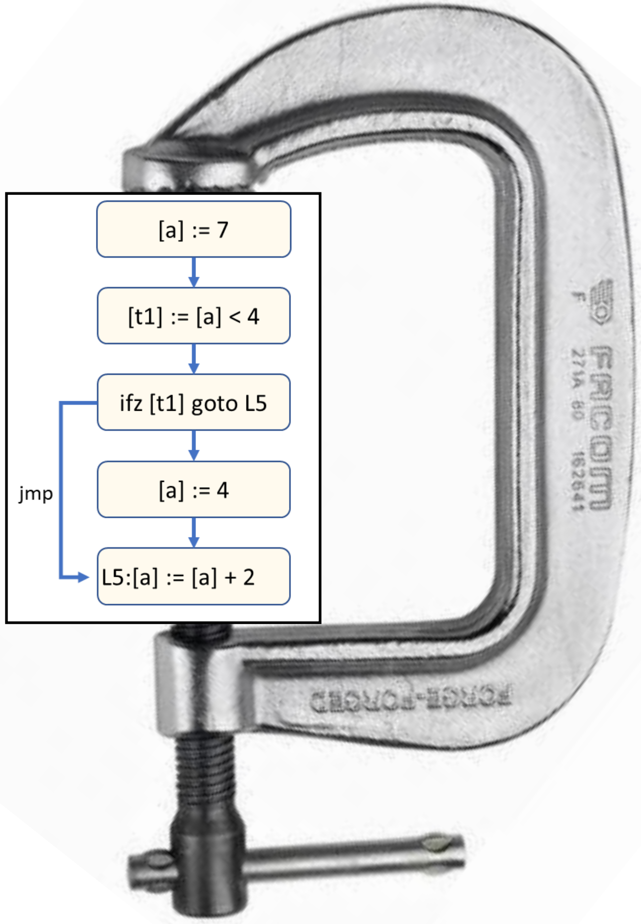
- Dead code elimination
- Common subexpression elimination
- Constant/copy propagation

Optimization

Intuition

Flowgraphs: Control flow graphs

- A more compact version of the instruction flow chart
- But still preserves the way in which control passes through the program

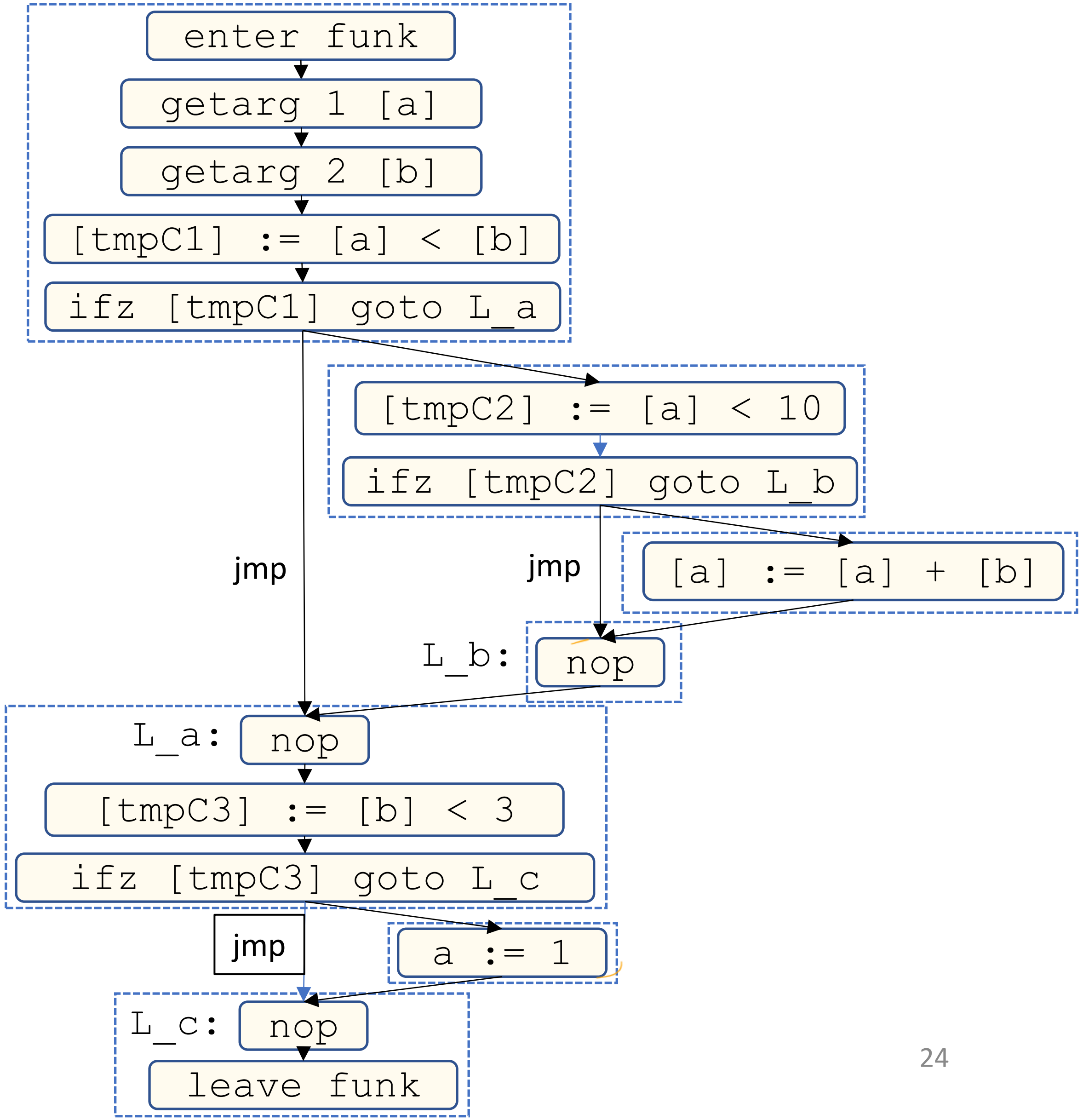


Compacting the Flowchart Concept

Flowgraphs: Control flow graphs

The flowchart is needlessly verbose

- We could put multiple instructions in a node
- Group the instructions that <u>always execute together</u>



Basic Blocks

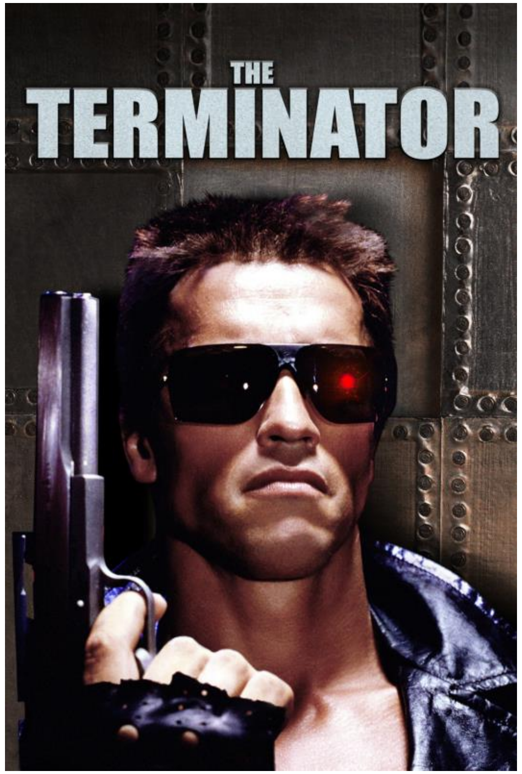
Flowgraphs: Control flow graphs

- Definition: Sequence of instructions guaranteed to execute without interruption

Basic Blocks Boundaries

Flowgraphs: Control flow graphs

- “Terminator” – An instruction that ends a basic block
- “Leader” – An instruction that begins a block



Basic Blocks

Flowgraphs: Control flow graphs

- Sequence of instructions guaranteed to execute without interruption
- Terminology:
  - "Leader" – An instruction that begins a block
    - *The first instruction in the procedure*
    - *The target of a jump*
    - *The instruction after an terminator*
  - "Terminator" – An instruction that ends a basic block
    - *The last instruction of the procedure*
    - *A jump (ifz, goto)*
    - *A call (We'll use a special LINK edge)*

Basic Blocks

Flowgraphs: Control flow graphs

Leaders

The first instruction in the procedure

The target of a jump

The instruction after a terminator

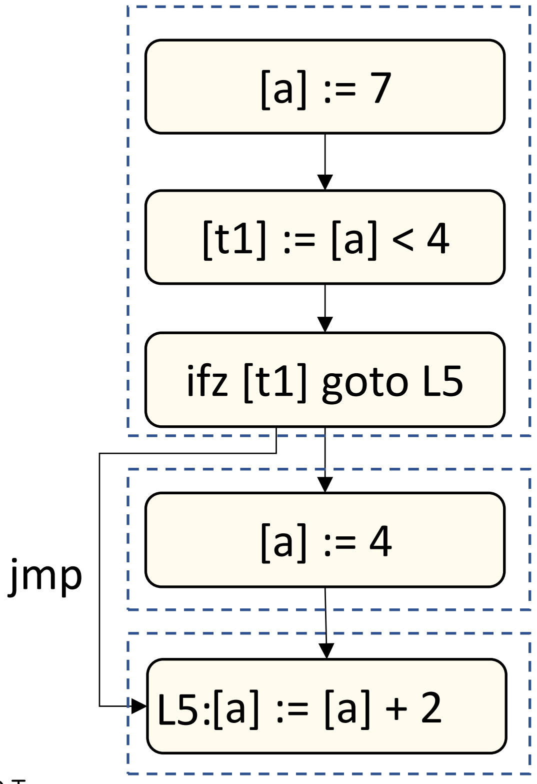
Terminators

The last instruction in the procedure

A jump (ifz, goto)

A call (We'll use a special LINK edge to successor)

Next instruction is a leader



Construction algorithm

```
foreach instr i in procedure:
   if i is a leader, begin a new BBL
   if i is a terminator, end current BBL
```

Building Basic Blocks

Flowgraphs: Control flow graphs

Leaders

The first instruction in the procedure

The target of a jump

The instruction after a terminator

Terminators

The last instruction in the procedure

A jump (ifz, goto)

A call (We'll use a special LINK edge to successor)
Next instruction is a leader

This algorithm isn't optimal, but we'll go with it

Construction algorithm

```
foreach instr i in procedure:
   if i is a leader, begin a new BBL
   if i is a terminator, end current BBL
```

<u>**example**</u>

```
    jmp L1
L1: nop
```

The Control Flow Graph: Summary

Flowgraphs: Control flow graphs

A graph of basic blocks

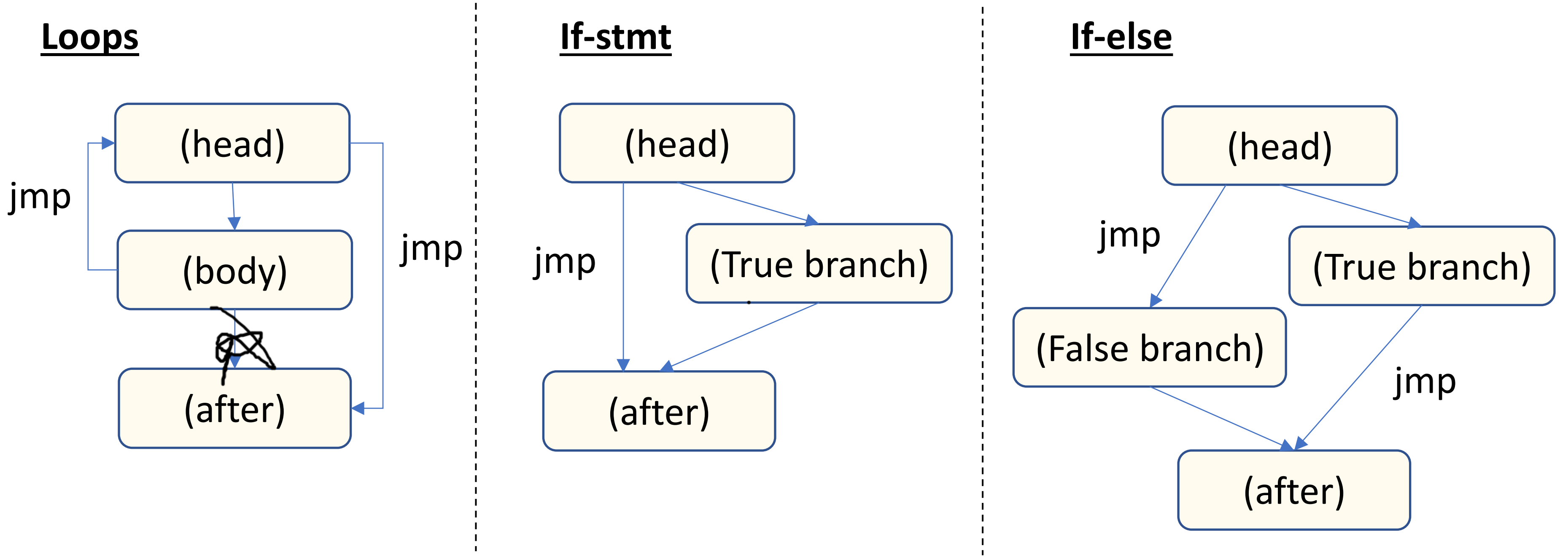
- One graph per procedure
  - Exactly one entry block
  - Exactly one exit block
- Distinguished edge types:
  - Back edges – an edge to a previously-encountered node
  - Call edge – Connects a call site to the called function
  - Link edge – Connects a function call to it's return point

Benefits of Basic Blocks

Flowgraphs: Control flow graphs

Makes CFGs a more manageable data structure

- Zoom out and observe procedure structure



Benefits of Basic Blocks

Makes CFGs a more manageable data structure

- Zoom out and observe procedure structure
- Zoom in to a BBL's "uninterrupted sequences"

Simplifies analysis:

- Many properties we want to know are trivial to compute within a BBL

Types of CFG Analysis

Control Flow Graphs: Representation

Modularizes analysis:

- Analysis within a single basic block

Traditionally called "Local" analysis

- Analysis between multiple basic blocks in a function

Traditionally called "Global" analysis

- What about analysis between multiple functions?

We'll come back to this one

Lecture Outline

Flowgraphs

Program analysis:

- Goals
- Control flow graphs

Local Optimizations

- Dead code elimination
- Common subexpression elimination
- Constant/copy propagation

Optimization

Dead Code Elimination

Flowgraphs: Local Optimizations

Remove "useless" instructions (those with no effect)

- Analysis: live variable analysis

This definition does not reach the end of the block nor any use in the block!

```
[x] := 1
[x] := 2
[y] := 3
[z] := [x] * [y]
write [z]
```

```
[x] := 2
[y] := 3
[z] := [x] * [y]
write [z]
```

Constant/Copy Propagation

Flowgraphs: Local Optimizations

Replace a variable use with its definition

- Analysis: “copy identification” (doesn’t really have a name)

Propagate x’s value

```
[y] := [x]
[z] := [x] * [y]
[y] := 1
write [z]
```

This assignment is now dead code!

```
[y] := [x]
[z] := [x] * [x]
[y] := 1
write [z]
```

When propagating constant values, can aid in constant folding

```
[x] := 1
[y] := 2
[z] := [x] + [y]
```

Fold to 3

```
[x] := 1
[y] := 2
[z] := 1 + 2
```

Common Subexpression Elimination

Flowgraphs: Local Optimizations

Reuse already-computed expressions

- Analysis: available expression analysis

```
[a] := [x] * [y]
[b] := [x] * [y]
[c] := [x] * [y]
```

⇒

```
[t] := [x] * [y]
[a] := [t]
[b] := [t]
[c] := [t]
```

Lecture End

Flowgraphs

Summary

- Control Flow Graphs serve as an abstraction of the routes through the program
- Basic blocks summarize guaranteed sequences and enable local optimizations (DCE, CP, CSE)

Next Time

- Global optimizations – extending optimization across multiple basic blocks