

ECS 665 COMPILER CONSTRUCTION

Postcompilation

Previously...

Review – The Heap

Heap Memory

- Using the heap
- OS interface

Garbage collection

- Reference Counting
- Mark and Sweep

You Should Know

- How to program with heap memory
- The basic concepts of the two garbage collection schemes



Machine Codegen



Today's Lecture

Postcompilation

Compiler Toolchains

- Overview

Component Definitions

- Detour - History
- What GCC Does

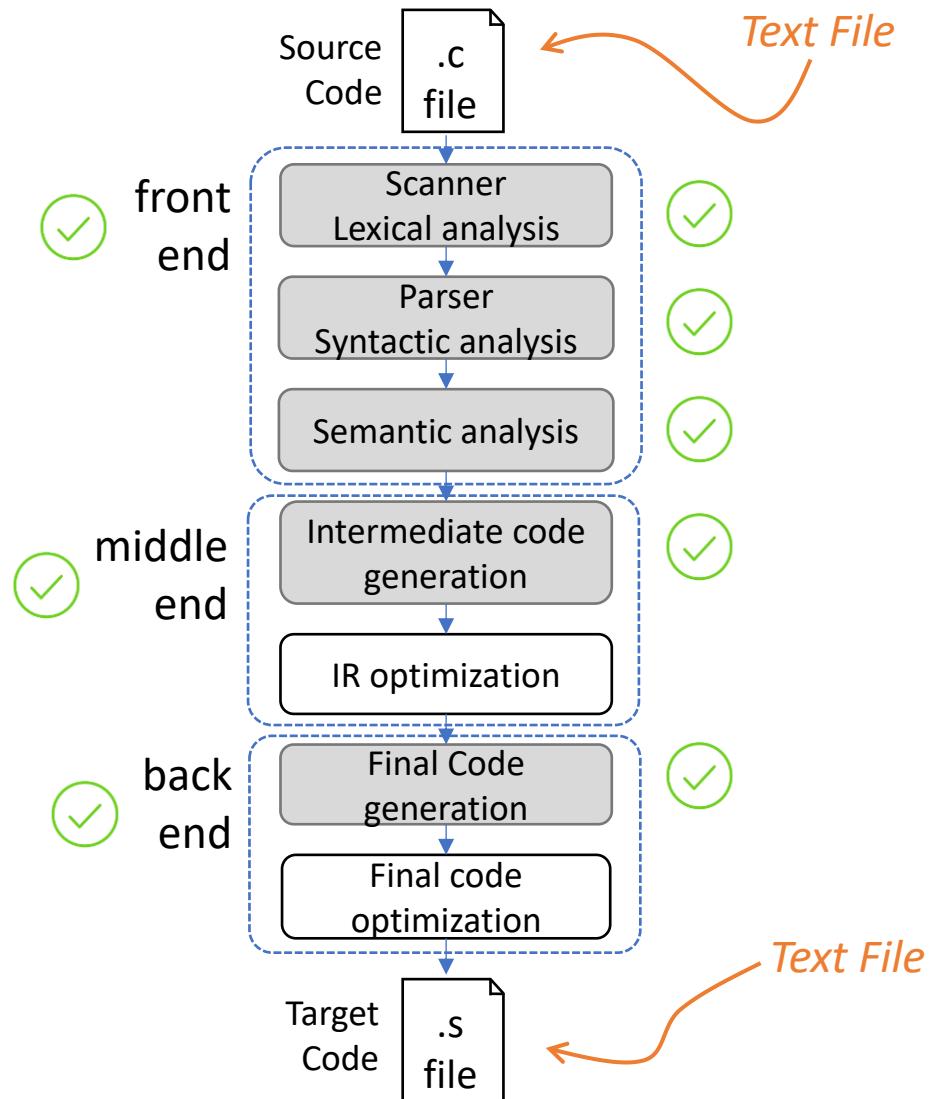
Component Walkthrough

- Assembler
- Linker
- Loader



Compiler Construction: Status

Progress Pics



Implemented all modules of a non-optimizing compiler

- Front end, middle end, and back end are all complete

One detail remains

- What use is an assembly text file?

OS Perspective: Still Just Text

Compiler Toolchains - Overview

From text file to text file

- Source code and assembly code equally useless to processor

Pretty unlike that runtime understands a text format

- Nearly always a better binary representation
- Some work left after the compiler has finished



*From text file to text file:
OS sees nothing accomplished*

OS Perspective: Still Just Text

Compiler Toolchains - Overview

From text file to text file

- Source code and assembly code equally useless to processor

Pretty unlike that runtime understands a text format

- Nearly always a better binary representation
- Some work left after the compiler has finished

“Shouldn’t the compiler target that?”



Compiler doesn't (need to)
output object code / executables

“But when I run gcc I get a binary”

- You, maybe

Then perhaps gcc is more than a compiler!



Did I just blow your mind?

Probably not, because it's just a matter of definitions

Compiler vs Compiler Toolchain

(A pedantic distinction)

Program lifecycle components often collectively referred to as “compiler” (*vs.* compiler toolchain)*

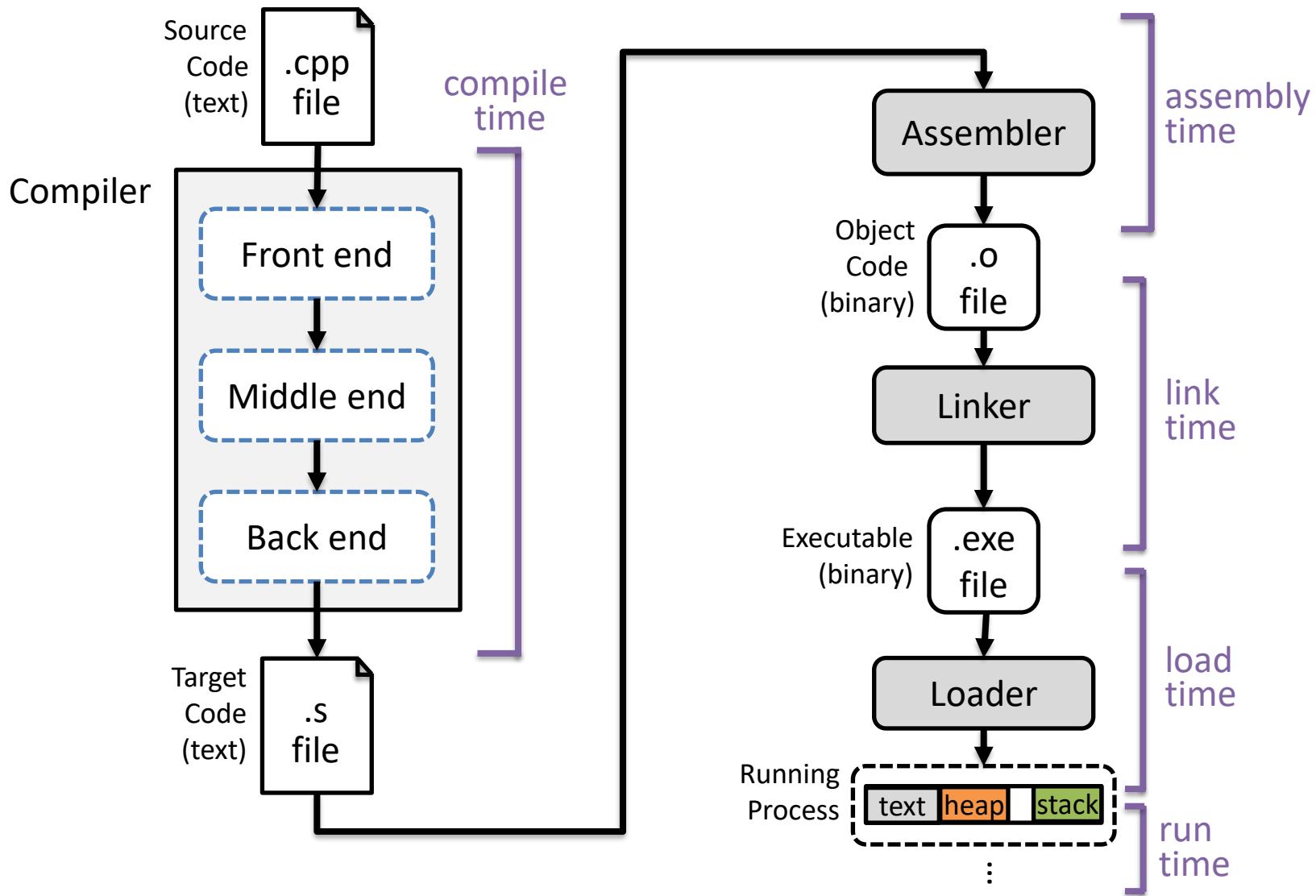
- These components *can* be invoked separately



[*]: Unfortunately (in my opinion) this convention kinda erases the term for the high-level text code to low-level text code component

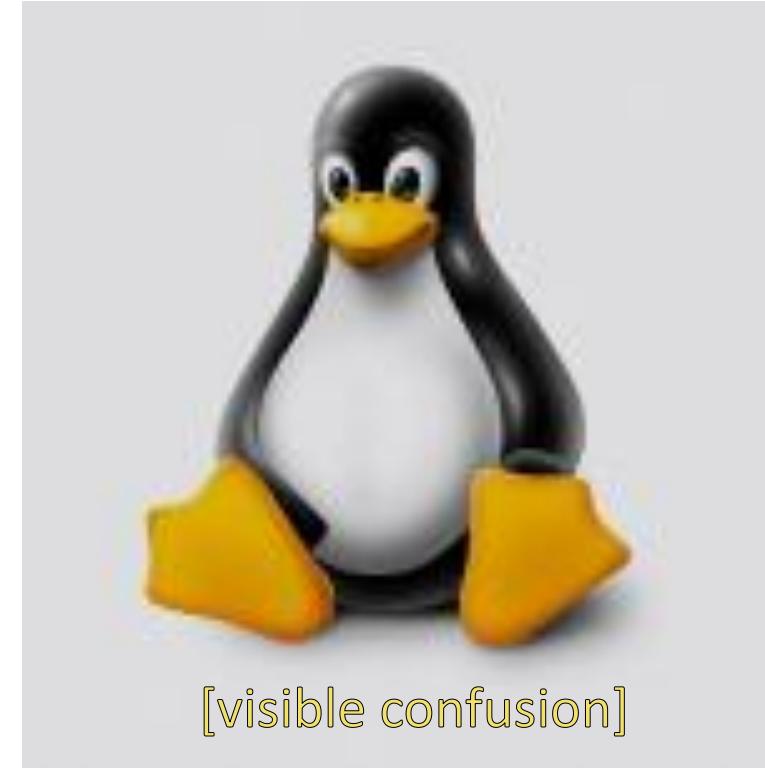
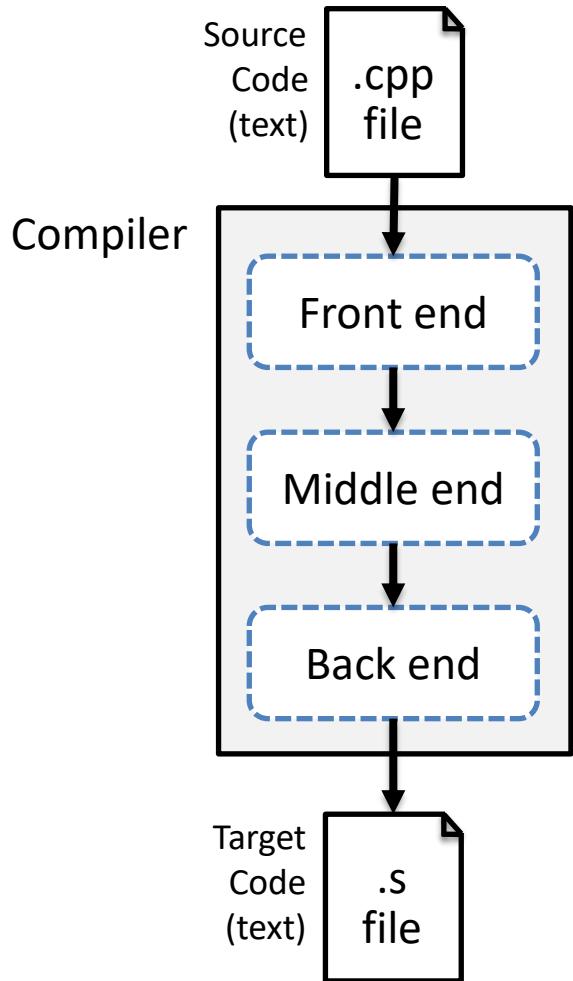
The Compiler Toolchain

Postcompilation - Component Walkthrough



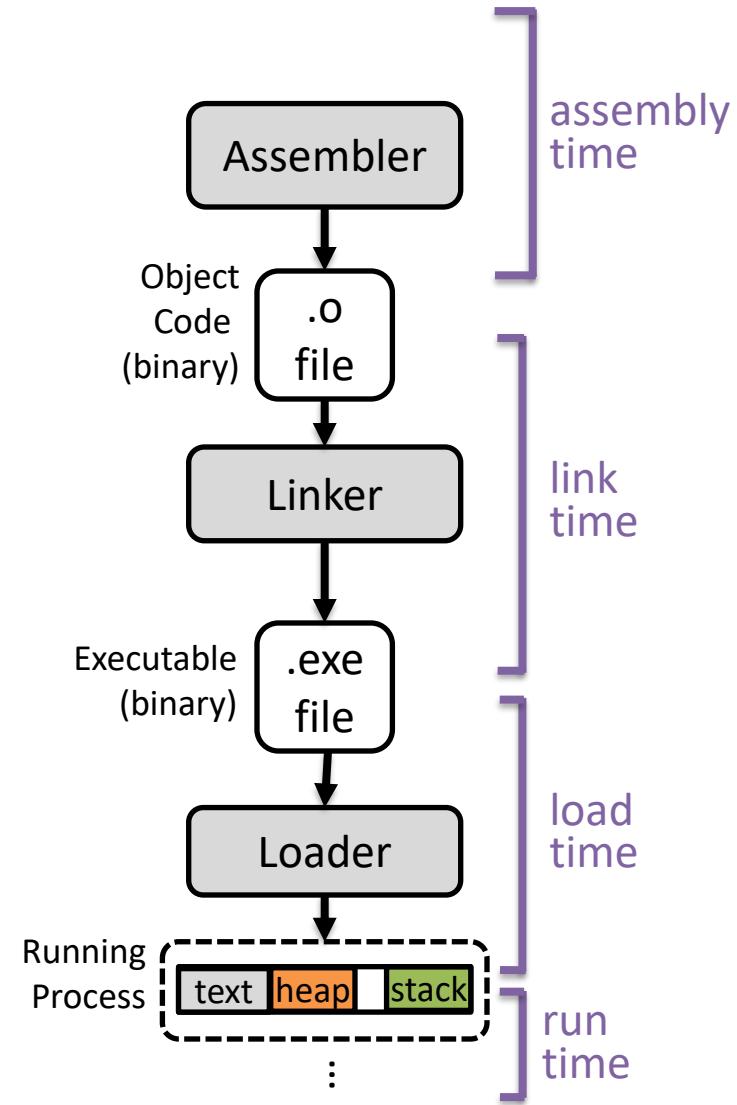
The Compiler Toolchain

Postcompilation - Component Walkthrough



The Compiler Toolchain

Postcompilation - Component Walkthrough



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Defining Toolchain Components

Postcompilation - Component Walkthrough

DEFINITIONS

(They're a little bit fuzzy)

Historical Factors

Detour: Component Terminology



- Compiler toolchain largely evolved from need to handle more complex programming tasks
 - Definitions have likewise evolved, with somewhat uneven consensus

Genesis of “Computer”

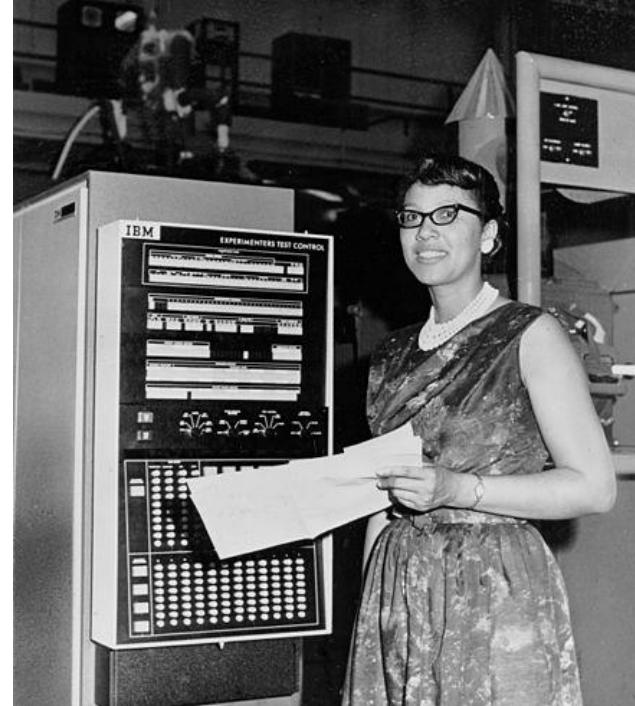
Detour: Component Terminology



From 1613-1945:

- “Computer” exclusively meant “a person who carries out computations”

Original programming consists of inputting numeric machine codes



Melba Roy Mouton: Early “human computer”

Genesis of “Assembler”

DETOUR

Detour: Component Terminology

1949: EDSAC features first code akin to modern notion of assembler

1953: Term “Assembler” popularized for a program that concretized symbolic programs



Maurice Wilkes: developed assembler for EDSAC



Nathaniel Rochester:
Published assembler
design

Genesis of “Linker”

Detour: Component Terminology

DETOUR

1952: A-0 system features first code akin to the modern notion of a linker (confusingly called a compiler at the time)

Allowed multiple programs to be combined together before being run



Grace Hopper: developed A-0

Genesis of “Compiler”

Detour: Component Terminology



1951: Publication of a translator from a high-level language to low-level code, considered the first compiler

(Also specifies the compiler in new high-level language)



Corrado Böhm: published first compiler in his PhD thesis

Genesis of “Loader”

Detour: Component Terminology

DETOUR

1947: ENIAC introduces the notion of relocating code to accommodate pre-defined subroutines, a feature somewhat akin to loaders (though pre-OS)



John Mauchley: Wrote about ENIAC procedures to relocate code in memory so that debugged subprograms could be loaded

Enough History for Now

Detour: Component Terminology



Practical upshot: terms have changed somewhat

Program Lifecycle for Our Purpose

Postcompilation - Component Walkthrough

- Compile high-level source code text to low-level assembly code text
 - Assemble assembly code into object code binary modules
 - Link object code into a single executable
 - Load program into memory and run
-
- The diagram illustrates the program lifecycle components. A vertical stack of four rounded rectangles represents the stages: Compiler (cc) at the top, followed by Assembler (as), Linker (ld), and Loader (loader) at the bottom. Brackets on the right side group these stages into two main sections: 'Compiler' (covering cc and as) and 'Linker' (covering ld and loader). A note to the right of the 'Assembler' stage states: 'These modules can be invoked separately!'
- Compiler
cc
- Assembler
as
- Linker
ld
- Loader
(launch program)
- These modules can be invoked separately!

Today's Lecture

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Component Walkthrough

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The gcc “subprograms”

Fun home experiment:

- Create empty.c

```
int main() { }
```

- Run gcc in verbose mode on empty.c

```
cycle1: gcc -v empty.c
```

```
a807d786@cycle1:~/libc_link$ gcc -v empty.c
Target: x86_64-linux-gnu
[...]
Thread model: posix
gcc version 7.5.0 (Ubuntu 7.5.0-3ubuntu1~18.04)
COLLECT_GCC_OPTIONS='-v' '-mtune=generic' '-march=x86-64'
/usr/lib/gcc/x86_64-linux-gnu/7/cc1 -quiet -v -imultiarch x86_64-linux-gnu empty.c -quiet -dumpbase empty.c -mtune=generic -march=x86-64 -auxbase empty -version -fstack-protector-strong -Wformat -Wformat-security -o /tmp/ccs6BRmG.s
[...]
GNU C11 (Ubuntu 7.5.0-3ubuntu1~18.04) version 7.5.0 (x86_64-linux-gnu)
  compiled by GNU C version 7.5.0, GMP version 6.1.2, MPFR version 4.0.1, MPC version 1.1.0, isl version isl-0.19-GMP
The compiler proper
Source code (input)  
  

The assembler
GGC heuristics: --param ggc-min-expand=100 --param ggc-min-heapsize=131072
Compiler executable checksum: b62ed4a2880cd4159476ea8293b72fa8
COLLECT_GCC_OPTIONS='-v' '-mtune=generic' '-march=x86-64'
as -v --64 -o /tmp/cc4FfCb1.o /tmp/ccs6BRmG.s
[...]
Assembly code
Object code
COLLECT_GCC_OPTIONS='-v' '-mtune=generic' '-march=x86-64'
/usr/lib/gcc/x86_64-linux-gnu/7/collect2 -plugin /usr/lib/gcc/x86_64-linux-gnu/7/liblto_plugin.so -plugin-opt=/usr/lib/gcc/x86_64-linux-gnu/7/lto-wrapper -plugin-opt=-fresolution=/tmp/ccAxjj1l.res -plugin-opt=-pass-through=-lgcc -plugin-opt=-pass-through=-lgcc_s -plugin-opt=-pass-through=-lc -plugin-opt=-pass-through=-lgcc -plugin-opt=-pass-through=-lgcc_s --build-id --eh-frame-hdr -m elf_x86_64 --hash-style=gnu --as-needed -dynamic-linker /lib64/ld-linux-x86-64.so.2 -pie -z now -z relro /usr/lib/gcc/x86_64-linux-gnu/7/..../..../x86_64-linux-gnu/Scrt1.o /usr/lib/gcc/x86_64-linux-gnu/7/..../..../x86_64-linux-gnu/crti.o /usr/lib/gcc/x86_64-linux-gnu/7/crtbeginS.o -L/usr/lib/gcc/x86_64-linux-gnu/7 -L/usr/lib/gcc/x86_64-linux-gnu/7/..../..../x86_64-linux-gnu -L/usr/lib/gcc/x86_64-linux-gnu/7/..../..../lib -L/lib/x86_64-linux-gnu -L/lib/..../lib -L/usr/lib/x86_64-linux-gnu -L/usr/lib/gcc/x86_64-linux-gnu/7/..../..> /tmp/cc4FfCb1.o -lgcc --push-state --as-needed -lgcc_s --pop-state -lc -lgcc --push-state --as-needed -lgcc_s --pop-state /usr/lib/gcc/x86_64-linux-gnu/7/crtn.o
COLLECT_GCC_OPTIONS='-v' '-mtune=generic' '-march=x86-64'  
  

Runtime components
The linker
```

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Component Walkthrough

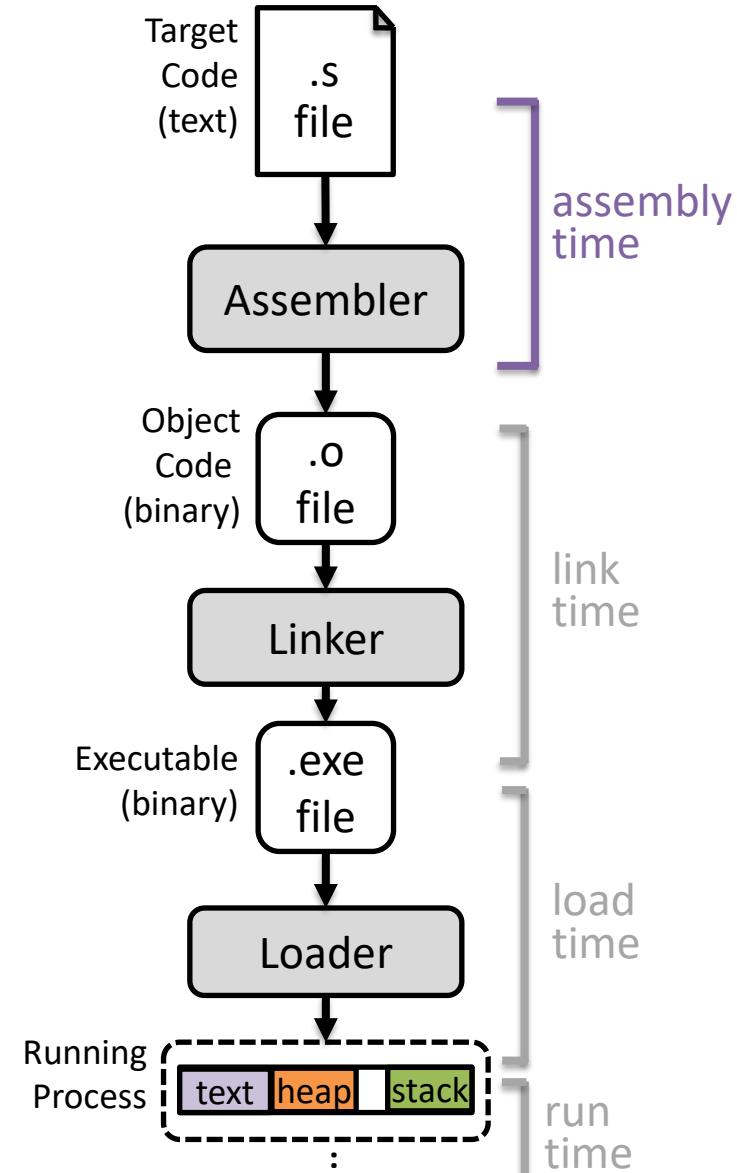
- Assembler
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The Assembler

Module Role and Operation

- Modern use:
 - Translates labels and instruction mnemonics to binary equivalent
- Output may still contain placeholders for some code and data



The Assembler: Responsibilities

Module Role and Operation

- Collect symbolic definitions (e.g. labels) into “low-level” symbol table
- Put code and data into segments
- Replace uses of labels with placeholder addresses (where available)

.data
_label1: .word

.text
jal _label7

.data
_str: .asciiz ...

.text
la \$t0 _label1

The diagram illustrates the assembly process. It starts with four separate code snippets arranged vertically. The top snippet is blue, the second is purple, the third is blue, and the bottom one is purple. An orange arrow points downwards from these snippets to a single orange box labeled "symbols". Below the "symbols" box are two stacked boxes: a blue one labeled ".data (binary)" and a purple one labeled ".text (binary)".

symbols

.data
(binary)

.text
(binary)

The Assembler

Module Role and Operation

src code

```
1 .data
2 gbl: .quad 12
3 .globl _start
4 .text
5 _start: jmp _start
6
```

```
as -o out.o file.s
objdump -Dwrt out.o
```

object code

```
1
2 code.o:      file format elf64-x86-64
3
4 SYMBOL TABLE:
5 0000000000000000 l    d  .text  0000000000000000 .text
6 0000000000000000 l    d  .data  0000000000000000 .data
7 0000000000000000 l    d  .bss   0000000000000000 .bss
8 0000000000000000 l    .data  0000000000000000 gbl
9 0000000000000000 g    .text  0000000000000000 _start
10
11
12
13 Disassembly of section .text:
14
15 0000000000000000 <_start>:
16 0: eb fe          jmp    0 <_start>
17
18 Disassembly of section .data:
19
20 0000000000000000 <gbl>:
21 0: 0c 00          or     $0x0,%al
22 2: 00 00          add    %al,(%rax)
23 4: 00 00          add    %al,(%rax)
24 ...
```

The Assembler: Invocation

Module Role and Operation

(to get assembly): cc -S ret3.c -o /tmp/a.s

as /tmp/a.s -o /tmp/a.o

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Component Walkthrough

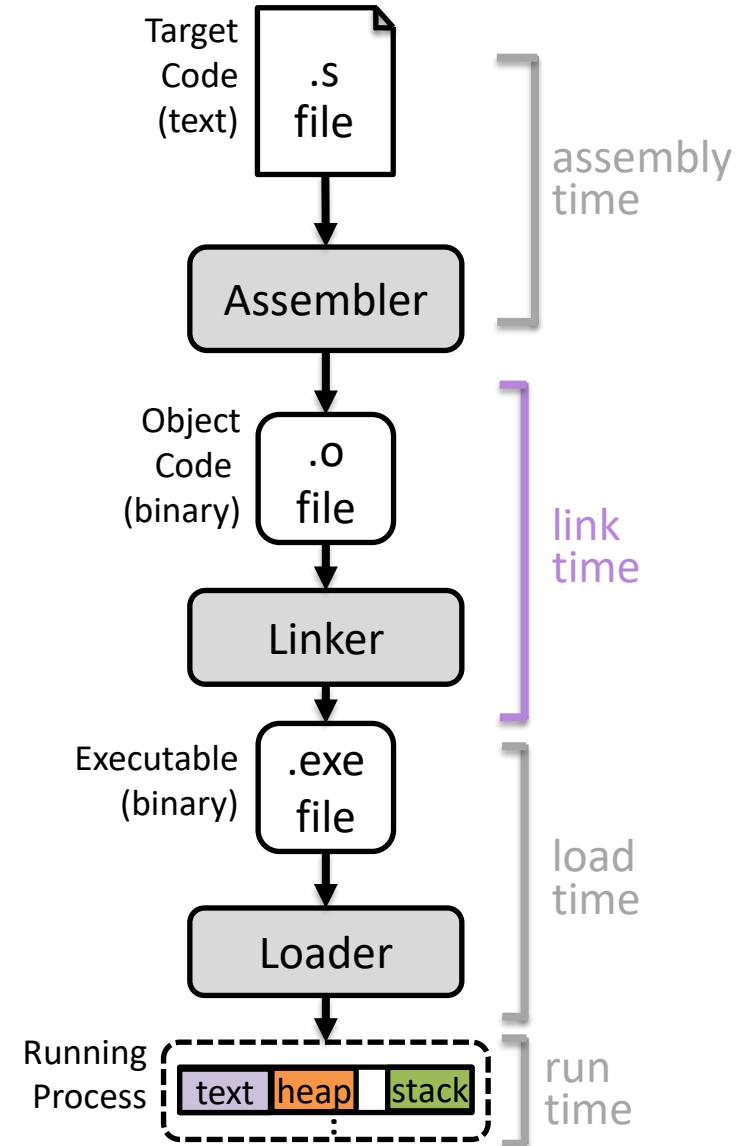
- Assembler
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The Linker

Module Role and Operation

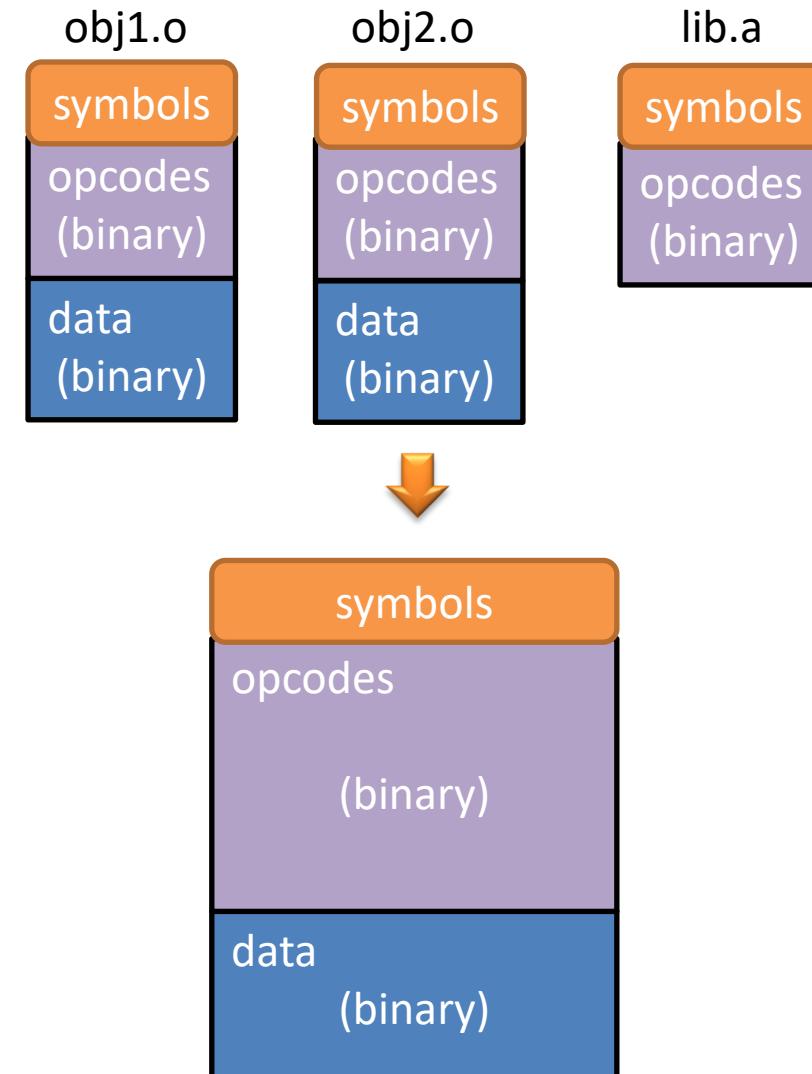
- Resolves placeholders betwixt multiple object files
- Stitches objects files (and static libraries) into a single binary



The Linker: Responsibilities

Module Role and Operation

- Calculate segment size
- Relocate segments
- Resolve symbol imports with exports where available
 - e.g. *call in obj1.o, callee in obj2.o*
- Add initialization routines
- Optimization!



The Linker

Module Role and Operation

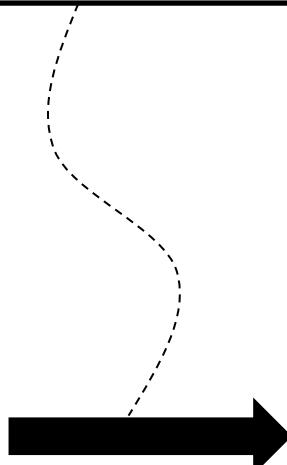
```
ld -o prog code.o  
objdump -Dwrt prog
```

Object code

```
1  
2 code.o:      file format elf64-x86-64  
3  
4 SYMBOL TABLE:  
5 0000000000000000 l    d  .text  0000000000000000 .text  
6 0000000000000000 l    d  .data  0000000000000000 .data  
7 0000000000000000 l    d  .bss   0000000000000000 .bss  
8 0000000000000000 l    .data  0000000000000000 gbl  
9 0000000000000000 g    .text   0000000000000000 _start  
10  
11  
12  
13 Disassembly of section .text:  
14  
15 0000000000000000 <_start>:  
16 0: eb fe          jmp    0 <_start>  
17  
18 Disassembly of section .data:  
19  
20 0000000000000000 <gbl>:  
21 0: 0c 00          or     $0x0,%al  
22 2: 00 00          add    %al,(%rax)  
23 4: 00 00          add    %al,(%rax)  
24 ...
```

Executable code

```
1  
2 code.prog:      file format elf64-x86-64  
3  
4 SYMBOL TABLE:  
5 0000000004000b0 l    d  .text  0000000000000000 .text  
6 0000000006000b2 l    d  .data  0000000000000000 .data  
7 0000000000000000 l    df *ABS* 0000000000000000 code.o  
8 0000000006000b2 l    .data  0000000000000000 gbl  
9 0000000004000b0 g    .text   0000000000000000 _start  
10 0000000006000ba g   .data  0000000000000000 __bss_start  
11 0000000006000ba g   .data  0000000000000000 __edata  
12 0000000006000c0 g   .data  0000000000000000 __end  
13  
14  
15  
16 Disassembly of section .text:  
17  
18 0000000004000b0 <_start>:  
19 4000b0: eb fe          jmp    4000b0 <_start>  
20  
21 Disassembly of section .data:  
22  
23 0000000006000b2 <gbl>:  
24 6000b2: 0c 00          or     $0x0,%al  
25 6000b4: 00 00          add    %al,(%rax)  
26 6000b6: 00 00          add    %al,(%rax)  
27 ...
```





The Linker: Simple Invocation

Module Role and Operation

The simplest use, “linking” a single object file:

```
ld prog.o -o prog.exe
```

More complex use, linking multiple object files

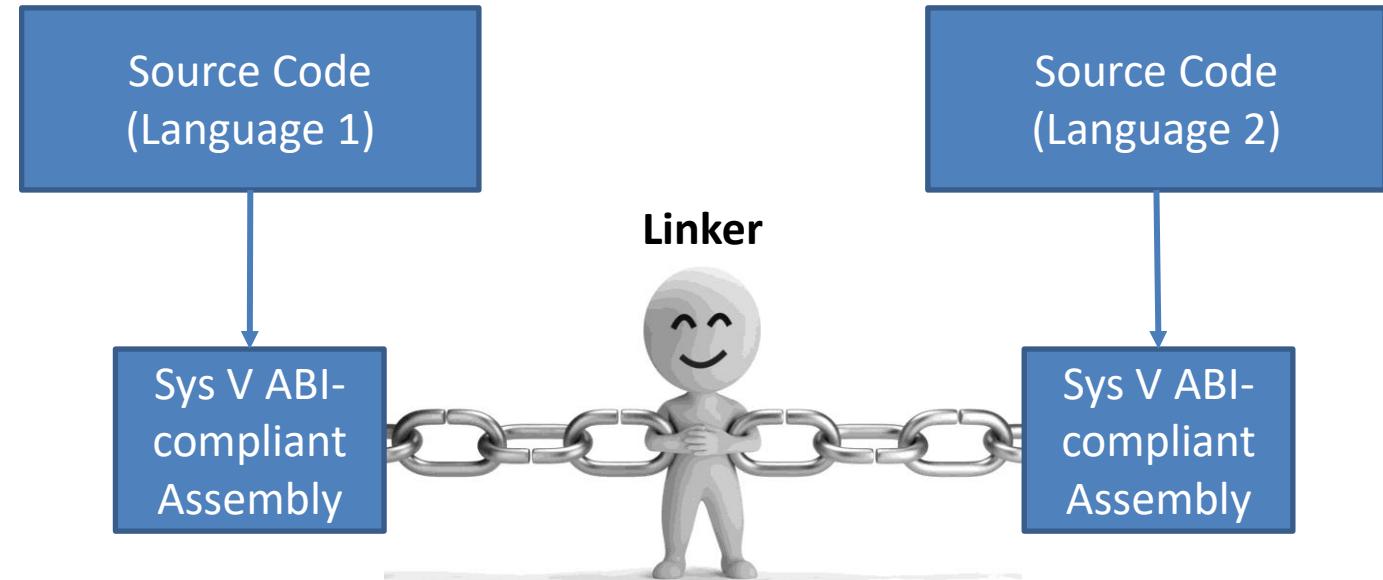
```
ld prog.o a.o b.o -o a.exe
```

Cross-language linking

Module Role and Operation

Object code is raw x64 code (containing placeholders)

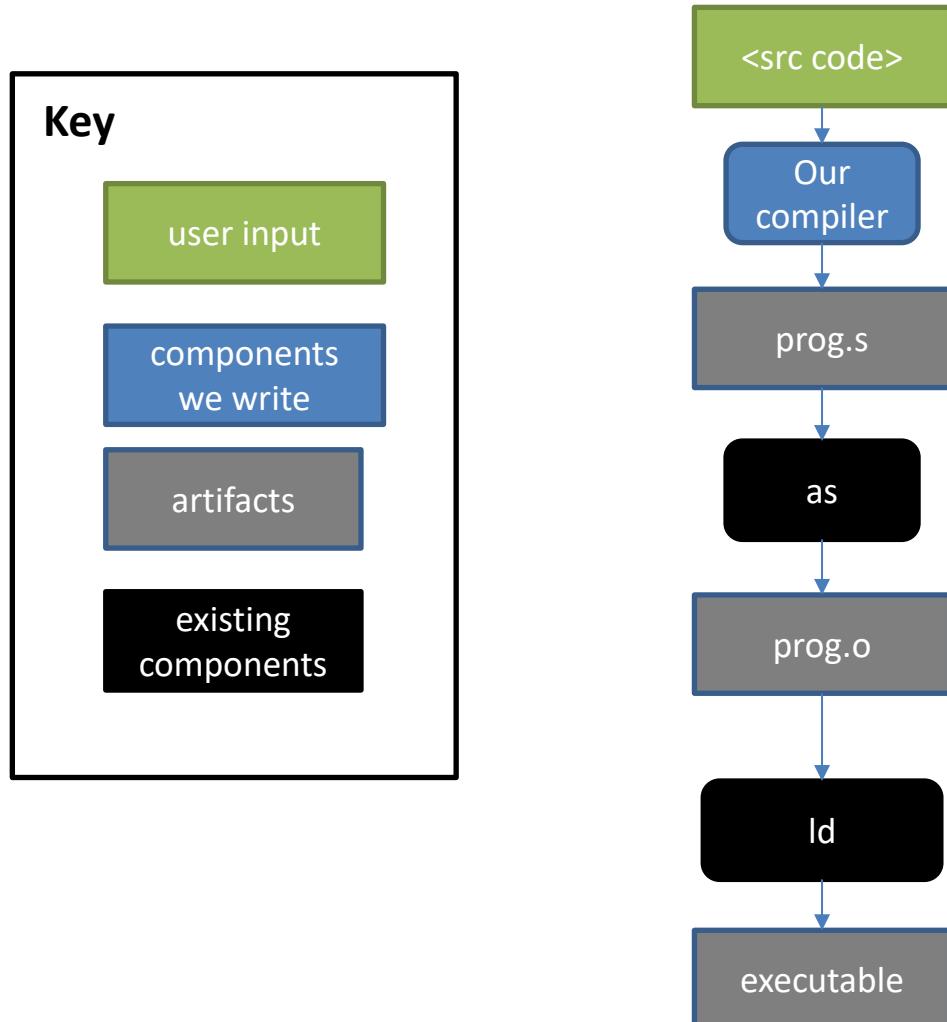
- Should obey the System V ABI
- We can link with object code from other languages!



Cross-language linking

Module Role and Operation

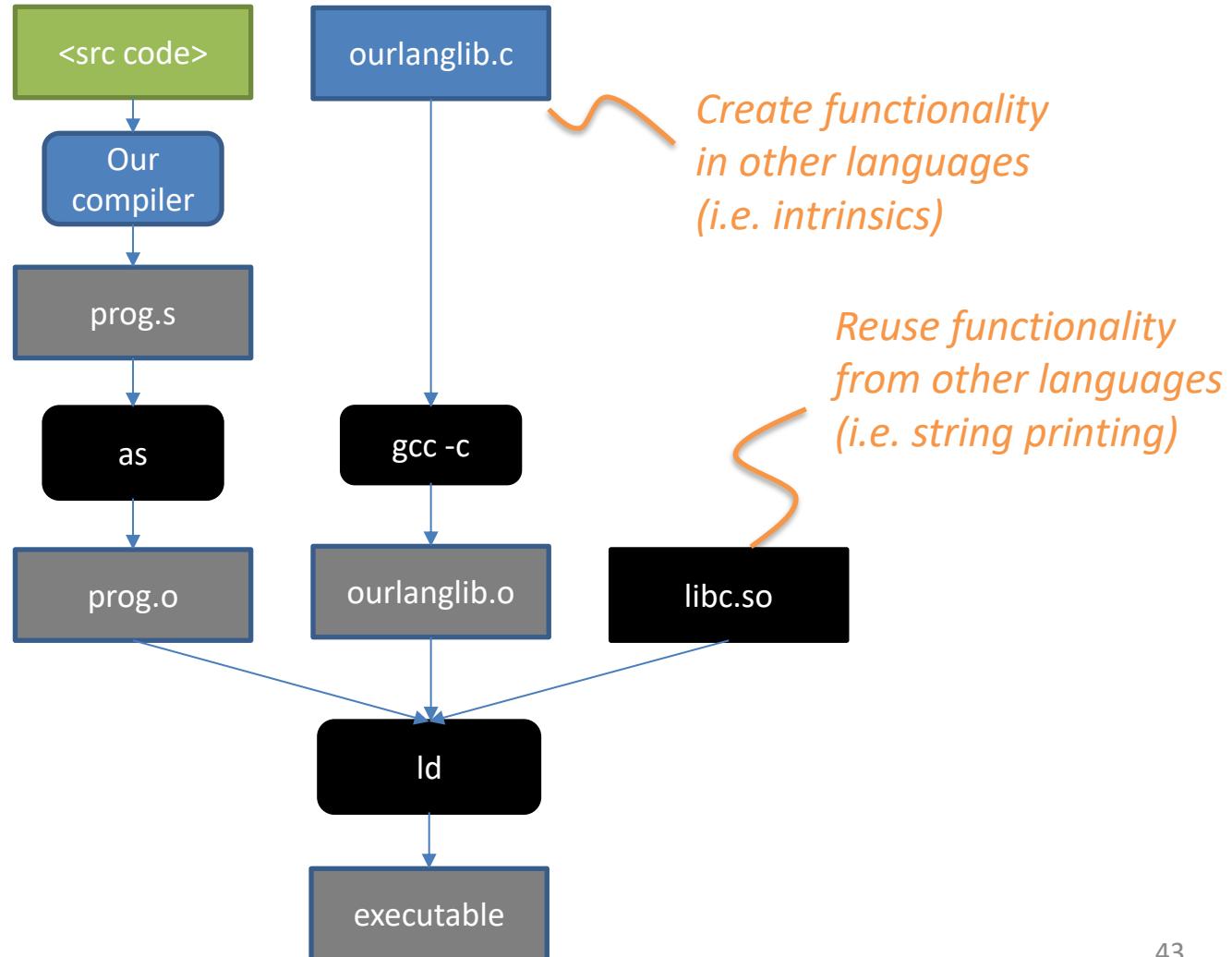
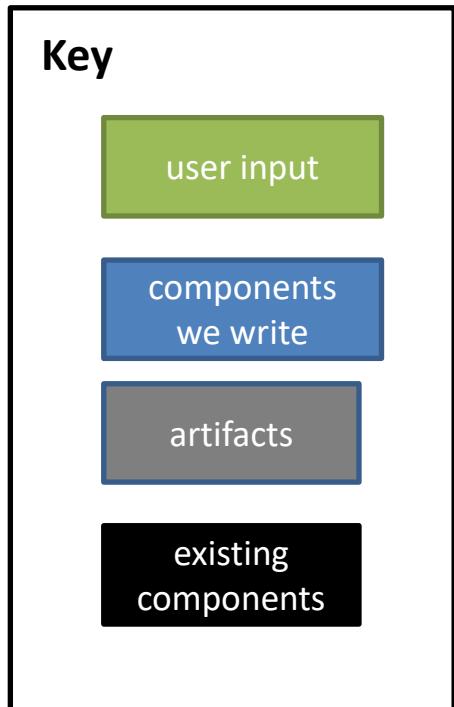
End-to-end compiler toolchain workflow



Cross-language linking

Module Role and Operation

End-to-end compiler toolchain workflow



Benefits of Cross-Linking

Module Role and Operation

Inter-language compatibility

- Let each language's strengths shine

Functionality reuse

- Avoid coding against the OS directly



Do not reinvent this



Why use the C Runtime?

Module Role and Operation

Src Code

```
int main() {  
    output 7;  
}
```

RYAN A. CHAPMAN



LATEST POST BROWSE POSTS

LINUX SYSTEM CALL TABLE FOR X86 64

PUBLISHED THU, NOV 29, 2012

https://blog.rchapman.org/posts/Linux_System_Call_Table_for_x86_64/

%rax	System call	%ordi	%rsi	%rdx
0	sys_read	unsigned int fd	char *buf	size_t count
1	sys_write	unsigned int fd	const char *buf	size_t count



Why use the C Runtime?

Module Role and Operation

Src Code

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int main() {  
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}
```

Asm Code

```
.globl _start  
.text  
_start: callq fn_main  
        movq %rax, %rdi  
        movq $60, %rax  
        syscall  
fn_main: <fn prologue>  

```

%rax	System call	%rdi	%rsi	%rdx
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Why use the C Runtime?

Module Role and Operation

Src Code

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int main() {  
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Asm Code

```
.data  
buffer: .asciz "?"  
.globl _start  
.text  
_start: callq fn_main  
        movq %rax, %rdi  
        movq $60, %rax  
        syscall  
fn_main: <fn prologue>  
        # output 7  
        movq $1, %rax #select sys_write  
        movq $1, %rdi #write to stdout  
        movq ???????, %rsi #buffer address  
        ...  
        syscall  
<fn epilogue>
```

%rax	System call	%rdi	%rsi	%rdx
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        movq $buffer, %rsi #buffer address  
        ...  
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fn_main: <fn prologue>  
        # output 7  
        movq $1, %rax #select sys_write  
        movq $1, %rdi #write to stdout  
        movq $buffer, %rsi #buffer address  
        movq $0x37, (buffer)#write to buffer  
        ...  
        syscall  
<fn epilogue>
```

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        movq $1, %rdi #write to stdout  
        movq $buffer, %rsi #buffer address  
        movq $0x37, (buffer)#write to buffer  
        movq $1, %rdx #string length  
        syscall  
<fn epilogue>
```

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        movq $1, %rdx #string length  
        syscall  
<fn epilogue>
```

It gets even
worse!



Why use the C Runtime?

Module Role and Operation

Src Code

```
...  
int main() {  
    output gbl;  
}
```

Asm Code

```
fn_main: <fn prologue>  
    # output gbl  
    movq $1, %rax #select sys_write  
    movq $1, %rdi #write to stdout  
    movq $buffer, %rsi #string to print  
    movq ??, (buffer) #write characters to buffer  
    movq ??, %rdx #string length  
    syscall  
<fn epilogue>
```

"I miss printf"

Prefill Buffer to all null bytes?

Loop over Buffer counting non-null bytes?

Repeatedly call idivq \$10?

Gbl: 173



Buffer:

0x31	0x37	0x33
ASCII 1	ASCII 7	ASCII 3



Use the C Runtime

Module Role and Operation

Src Code

```
int main() {  
    output 7;  
}
```

Old assembly

```
.globl _start  
.data  
buffer: .asciz "?"  
.text  
_start: callq fn_main  
        movq %rax, %rdi  
        movq $60, %rax  
        syscall  
  
fn_main: <fn prologue>  
        # output 7  
        movq $1, %rax #select sys_write  
        movq $1, %rdi #write to stdout  
        movq $buffer, %rsi #string to print  
        movq $31, (buffer) #write to buffer  
        movq $1, %rdx #string length  
        syscall  
<fn epilogue>
```

New Assembly

```
.text  
main: <fn prologue>  
      # output 7  
      movq $7, %rdi #int to print  
      callq myPrintInt  
<fn epilogue>
```

myStdLib

```
#include <stdio.h>  
#include <inttypes.h>  
void myPrintInt(int64_t arg){  
    fprintf(stdout, "%ld", arg);  
}
```



Use the C Runtime

Module Role and Operation

Src Code

```
int main() {  
    output 7;  
}
```

prog.o

User code
(calls myPrintInt)

mystdlib.o

Shim code
(calls fprintf)

C runtime

fprintf definition
_start label
and so much more!

New Assembly

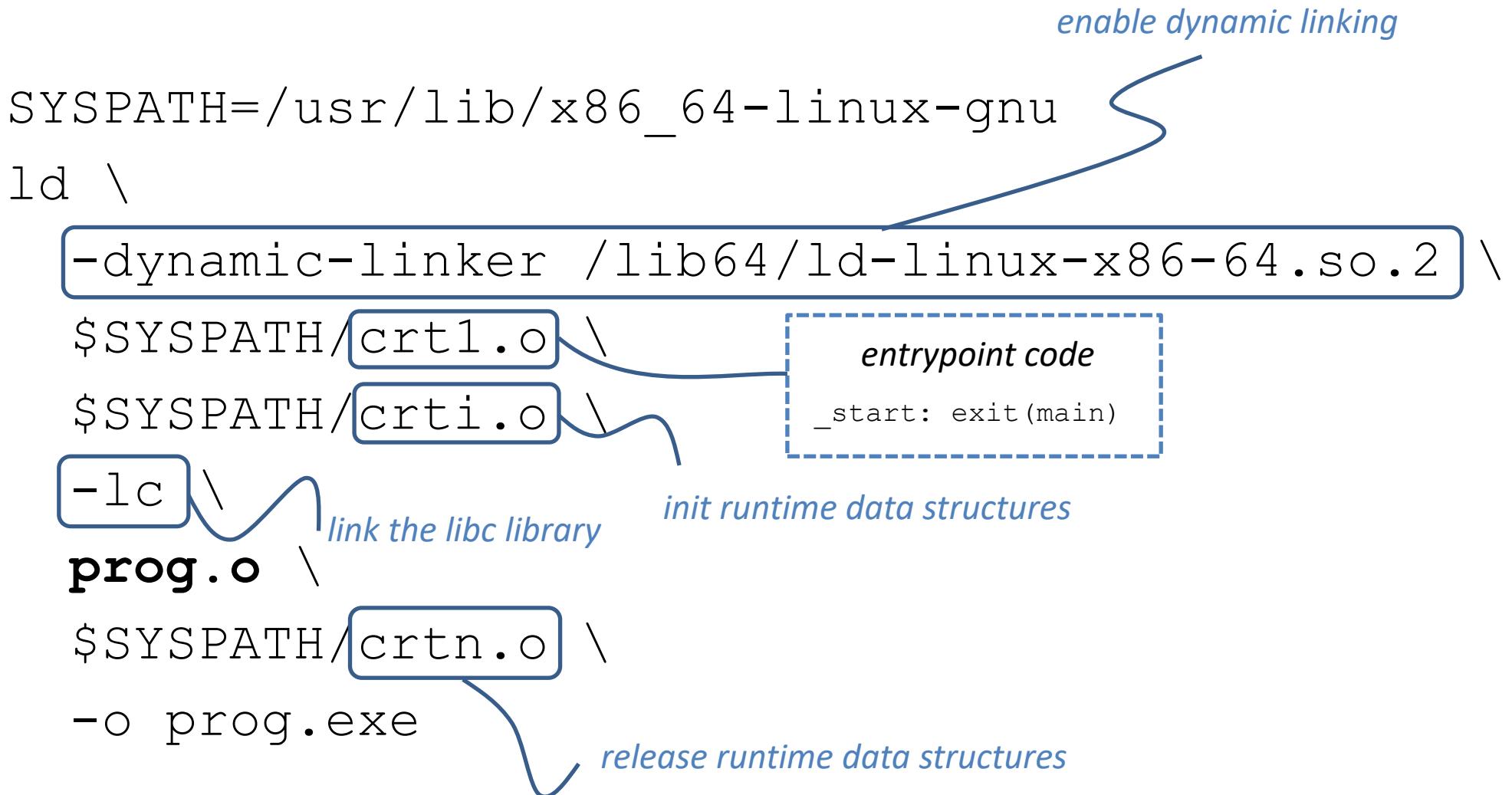
```
.text  
main: <fn prologue>  
        # output 7  
        movq $7, %rdi #int to print  
        callq myPrintInt  
<fn epilogue>
```

myStdLib

```
#include <stdio.h>  
#include <inttypes.h>  
void myPrintInt(int64_t arg) {  
    fprintf(stdout, "%ld", arg);  
}
```



Basing our runtime on C's Module Role and Operation



Today's Lecture

Postcompilation

Compiler Toolchains

- Overview

Component Definitions

- Detour - History
- What GCC Does

Component Walkthrough

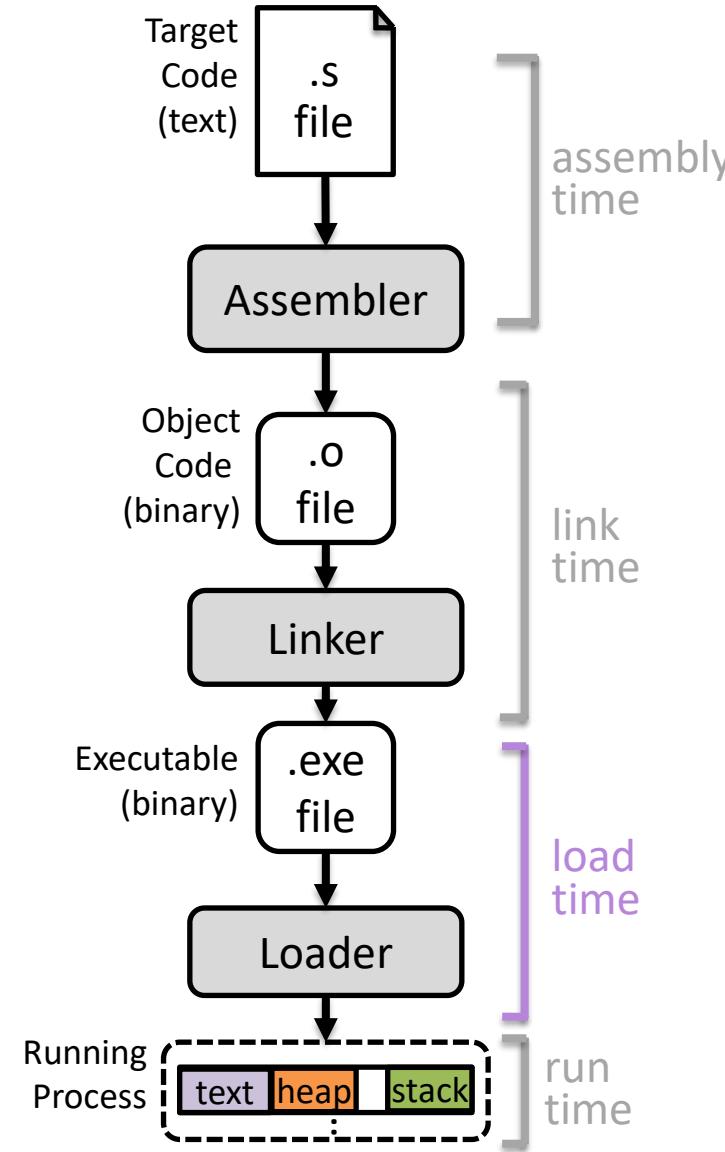
- Assembler
- Linker
- Loader



The Loader

Module Role and Operation

- Move code from secondary storage (disk) into primary storage (RAM)
- Resolves placeholders for dynamically-linked libraries (e.g. .so files)
- Kick off the program



The Loader

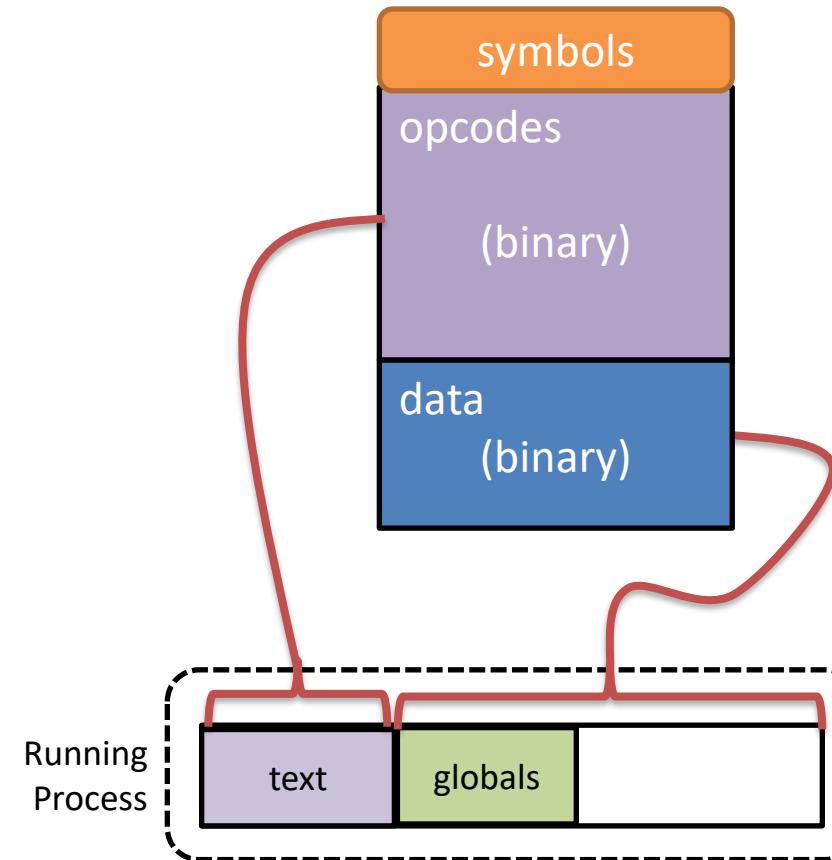
Module Role and Operation

Historically:

- Map program code into free memory
- May need to relocate symbols again

Modern Loaders:

- Simpler, as virtual memory handles physical address mapping
- Resolves dynamically-loaded libraries (.so, .dll)
- Part of the Operating System



That's all the Components!

Postcompilation: Underview

This basic process completes the “lifecycle” from source code to process

- From source to assembly
- From assembly to object code
- From object code to executable
- From executable to running process

Lecture Summary

Postcompilation: Underview

This Lecture

- Conceptual workflow from source text to process
- Tooling to perform this process
- Using cross-language linking to avoid re-implementation



Next up: Optimization!

- Improving on the naïve code generation we've performed
 - Better machine code
 - Better intermediate code