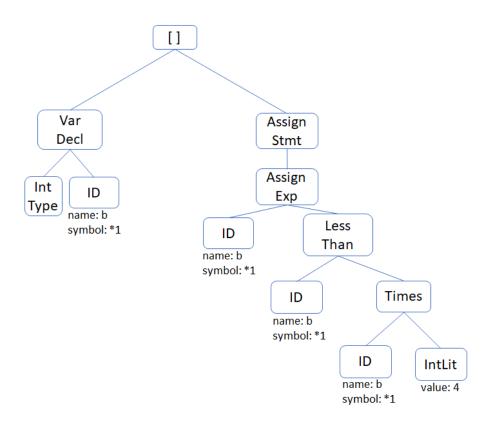
Check-In Review – Type Checking

Assume a program snippet has generated the following AST. Annotate each node with the type it corresponds to (or error if it is an error type). If a type analysis would issue a report, indicate that as well.





Check-in

Accidentally listed as a freebee?

Exam 2

Friday of next week

7:00 on Wednesdy of next week, roum TBA

University of Kansas | Drew Davidson

COMPLER CONSTRUCTION

Error Reporting

Last Time

Lecture Review – Type Analysis

Types

- What they are
- Why we have them

Type Rules

Examples

Connecting operations to their types

Enrich our static analysis pass

You Should Know

- The meaning of different aspects of type systems
- The simple AST-based type analysis
- How to propagate type errors



Semantics

Handling Errors

Type Analysis – Implementing Type Checking

- We'd like all distinct errors at the same time
 - Don't give up at the first error
 - Don't report the same error multiple times
- When you get error as an operand
 - Don't (re)report an error
 - Again, pass error up the tree



Type Error Example

Type Analysis – Implementing Type Checking int a; StmtList bool b; = true + 1 + 2 + b; AssignStmt AssignStmt b = 2;error **REPORT** error AssignExp AssignExp int bool error int Plus **REPORT** IdNode symbol IdNode IntLit bool error IdNode Plus type: int name: a int error IntLit type: bool **REPORT** Plus name: b int bool IntLit **BoolLit** 6 true

Today's Outline Lecture Overview – Error Reporting

Error Checking

- What counts as a bad program?
- How do we detect bad programs?

Limits of Analysis

The halting problem

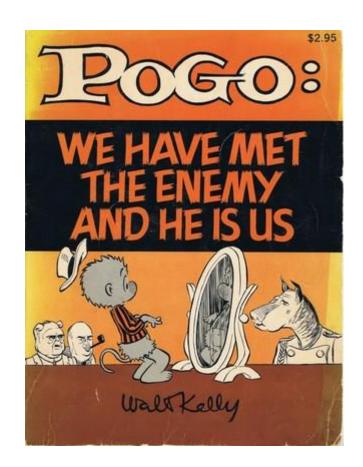


Semantics

Error Checking Semantic Analysis

Goal: save programmers from themselves

- It's not enough to compile the programmer's code
- Need to figure out what programmer meant to code



Quick Audience Poll

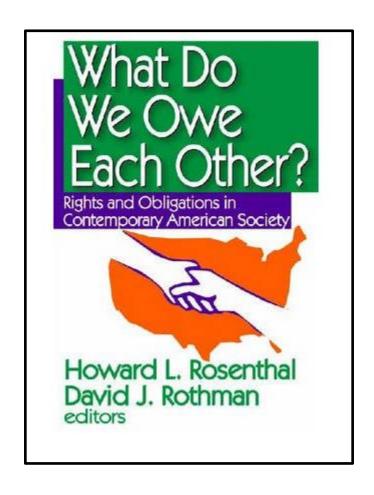
Semantics – Error Checking

Does this C program compile?

Should this C code compile?

```
int a = 0;
int main() {
   if (0 == 1) {
      b = 6;
   }
   return a;
}
```

A Compiler's Error-Checking Obligation Semantics - Error Checking



Understandability / Consistency

Compiler As Mind Reader Semantic Analysis – Broad View



A machine that infers your intent

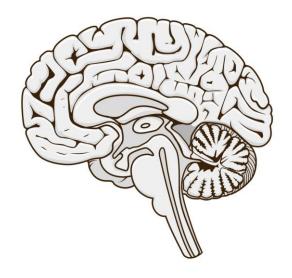
Compiler as Complainer Semantic Analysis – Broad View

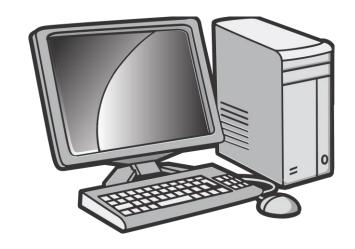


A grumpy old man that yells at you for breaking the rules

The Compiler Before the Compiler

Semantic Analysis – Broad View





Semantic gap: difference between the description of the same object in two different representation

Bug Hunting Semantic Analysis – Broad View

How do we prevent nonsense code from executing?

- We'll consider two ways of analysis:
 - Static
 - Dynamic



Putting guardrails on computation

Compiler Perspective

Semantic Analysis – Broad View

Static

Code analysis without execution

Dynamic

 Code analysis through execution Checks done at compile time

Analysis part of the compiler

itself

Checks done at run time

Analysis embedded into the program

Compiler Focus: Static Analysis Semantic Analysis – Broad View

Doesn't slow the program down

- Ok to take longer
- Ok to apply more heavyweight analysis

Has a "holistic" view of the program

- Has access to source code
- Knowledge of non-executed program paths

Limits of Error Checking Static Analysis

We'd LOVE to ensure bugfree programs

 Observe and report bugs before they are encountered

Usually we can't do this

Limits of static analysis



Limits of Static Analysis Static Analysis

Theoretical argument



Practical argument



The Halting Problem Static Analysis

Does a computation ever terminate?

Given a description of a Turing machine and its initial input, determine whether the program, when executed on this input, ever halts (completes). The alternative is that it runs forever without halting



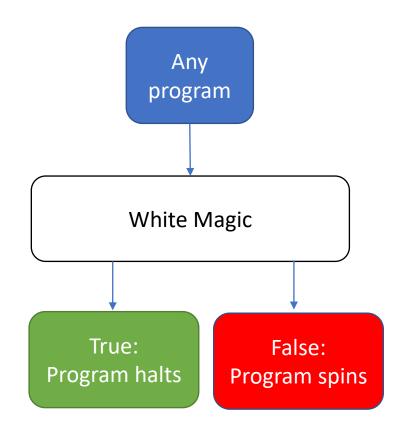
Sketching the Halting Problem Static Analysis

Effective procedure

 a procedure that is always yields a correct result on any input

Effective method for the halting problem would say:

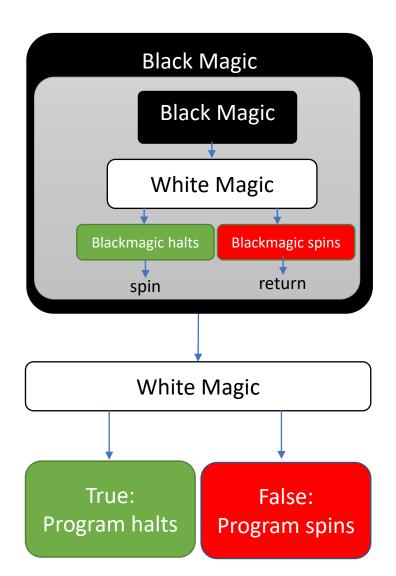
Return "true" if the program halts on the given input Return "false" if the program never halts on the given input



No Effective Method for Halting Static Analysis

assume white_magic(Function p)
returns true if p halts, false if p does not

```
void black_magic() {
    if white_magic(black_magic) {
        while true { }
    }
}
```



Implications of the Halting Problem Static Analysis

What does this have to do with, say, a null pointer analysis?

 No halting solution means no reachability solution

```
int * a = nullptr;
int main() {
    if (a != nullptr) {
        *a = 1;
    }
    return a;
}
```

Rice's Theorem Static Analysis

"All non-trivial semantic properties of programs are undecidable"



Rice's Theorem — Basic Idea

Static Analysis – Limits of Error Checking

What does this have to do with, say, a null pointer analysis?

No halting means no reachability

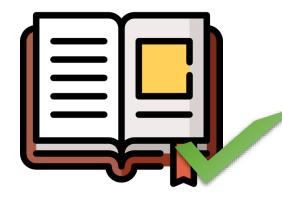
```
int main() {
   if (black_magic()) {
      int * p = 0;
      *p = 42;
   } else {
      return 0;
   }
}
```

Rice's Theorem - Implications Static Analysis - Limits of Error Checking

- We'd like to perfectly capture all bugs
 - We can't be right all of the time
 - We <u>can</u> choose **HOW** we are wrong

Limits of Static Analysis Static Analysis

Theoretical argument



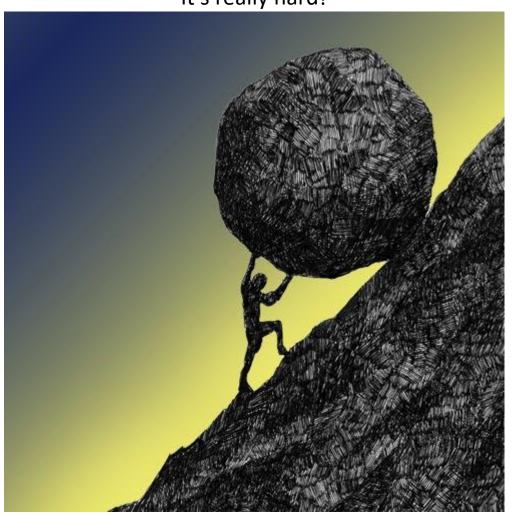
Practical argument



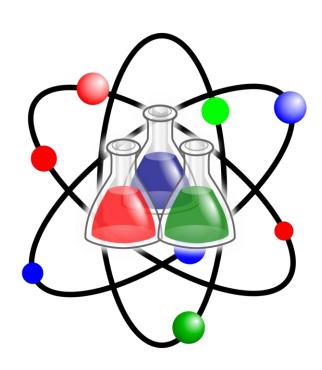
What if we only consider the universe of programs not written by (bleep) heads?

Practical Argument Static Analysis

It's really hard!

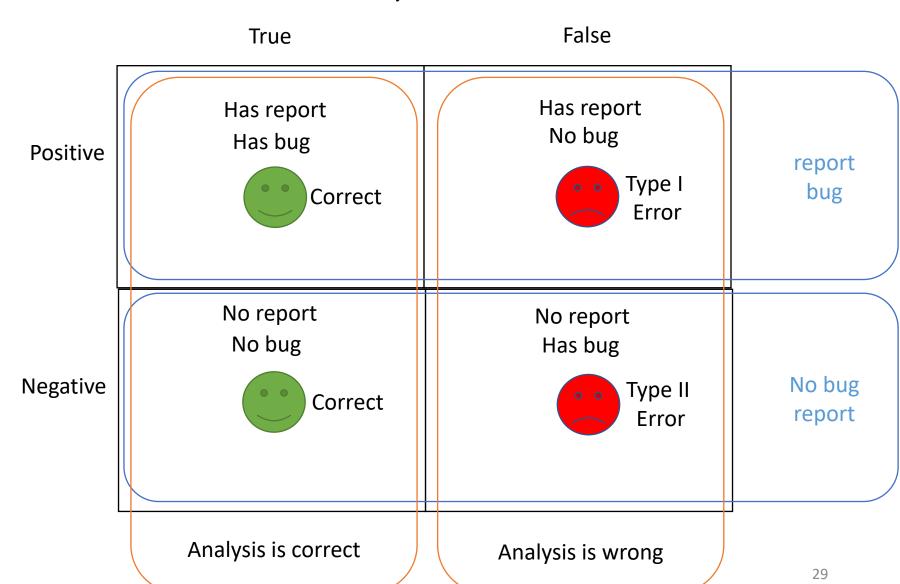


Let's do some Sciency-Sounding Stuff Static Analysis - Evaluation



Evaluating a Bug Detector

Static Analysis - Evaluation



Guarantees Under Imperfect Detection Static Analysis – Limits of Error Checking

Consistency / Reliability super important for users

We'd like to limit the <u>kinds</u> of errors we report

We can choose which type of bug report error to avoid

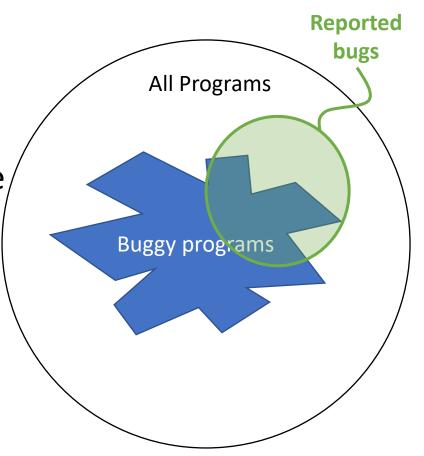
- Soundness: No false positives
- Completeness: No false negatives

VISUAL ANALOGY Static Analysis – Limits of Error Checking

Imagine the universe of all programs is contained in a circle

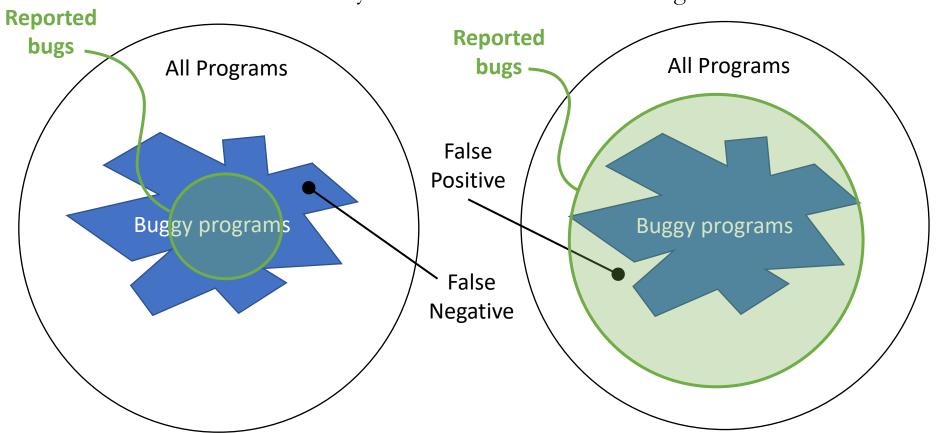
 You can draw a circle around the programs you report as buggy

 The actual buggy programs occupy a jagged region



Soundness and Completeness

Static Analysis – Limits of Error Checking



Sound bug detection

All correct programs pass through (No false positive problem)

Some buggy programs pass through (has false negative problem)

Complete bug detection

All buggy programs get flagged (No false negative problem)

Some correct programs get flagged (has false positive problem)

Partial Correctness Static Analysis – Limits of Error Checking

- Make best-effort procedures that are neither sound nor complete
- We can analyze the result of a statement under certain assumptions
 - Assume that the statement is executed
 - Assume that the statement actually completes