

Check-In

Scope Review

Does this program compile in a static scoping scheme? In a dynamic scoping scheme?
What is the output for each scheme in which it compiles?

```
a : int;
v : () void {
    a = a + 1;
    give a;
}
w : () void {
    a = 7;
    v();
    give a;
}
main : () void {
    int a;
    w();
    give a;
}
```

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ECCS 665
COMPILER
CONSTRUCTION

The title is presented in a highly stylized, futuristic font. The words "COMPILER" and "CONSTRUCTION" are in a bold, metallic, 3D style with a blue and white glow. "ECCS 665" is written in a pink, cursive script above "COMPILER". The background is a dark blue space with stars and a glowing purple grid floor.

Semantic Analysis

Announcements

Administrivia

Announcements

Administrivia

P3 Questions

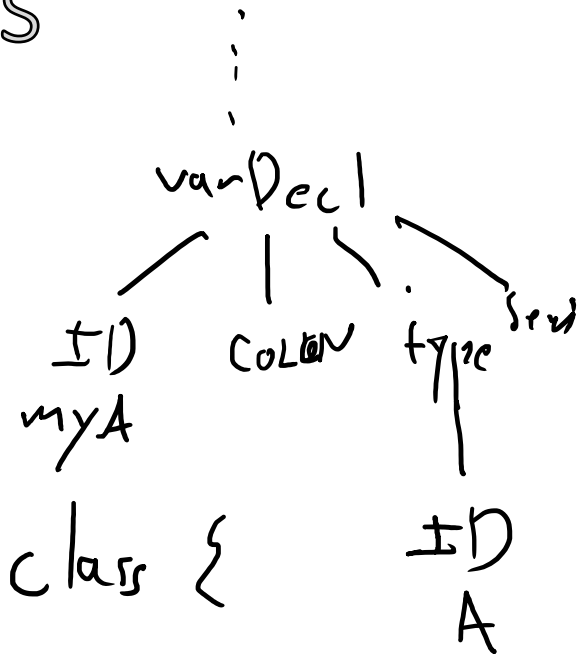
- ### P3 Questions
- ✓ - The POSTDec operator
 - - The MAGIC operator
 - Class-type variables and declarations

```
class A {
    int b
```

5.

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```
int main() {
    A a;
    a.b = 2;
}
```



```
A: class {  
    b: int;
```

§

1
2
3

```
main : ( ) int {  
    my A : A;  
    a -- b = 2;  
}
```

Last Lecture

Lecture Review – Scope

Issues of Scope

- Scheme
- Shadowing
- Overloading

You Should Know

- Scope properties
- How scope affects semantics
- High-level scope rules for our language

```
int a;  
if (true) {  
    int a;  
    ...  
    a = 4  
}
```



Lecture Outline

Lecture Overview – Semantic Analysis

Name Analysis

- Enforcing scope

Symbol Table

- What it is
- What it does



Semantics

Name Analysis

Semantic Analysis

Idea:

- Associates IDs with their uses in the program
 - i.e. Emplace symbol table entries
- Implemented as an AST pass

Purpose:

- Needed for code generation
- Catch some obvious errors
 - (undeclared IDs)



Recognizing Identifier Context

Semantic Analysis

A Context-Free Grammar is... context free

- A node needs information from outside of its subtree
- The definition of a identifiers needs to be connected to its occurrences
- We need a data structure to propagate such context



There's lots of ways to do this!
We'll just cover 1 way

The Compiler's Symbol Table

Semantic Analysis

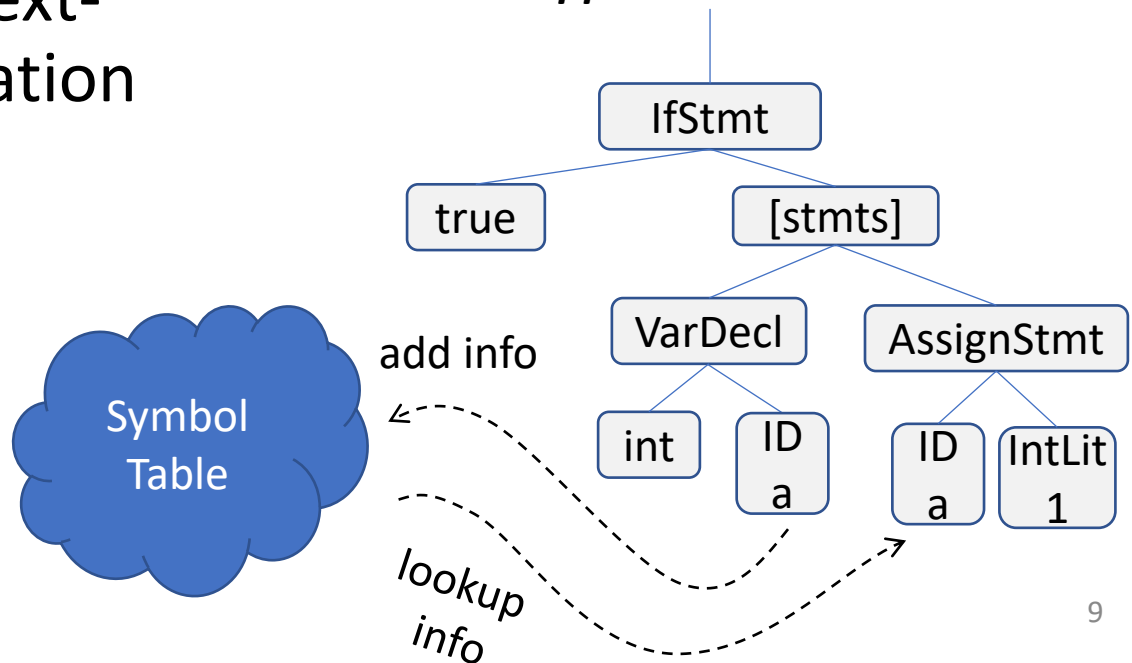
Repository of semantic symbol information

- Populated during a walk of the AST
- Propagates context-sensitive information

Program Snippet

```
if(true){  
    int a;  
    a = 1;  
}
```

AST Snippet



The Compiler's Symbol Table

Semantic Analysis

What's in the symbol table:

- Depends on the language

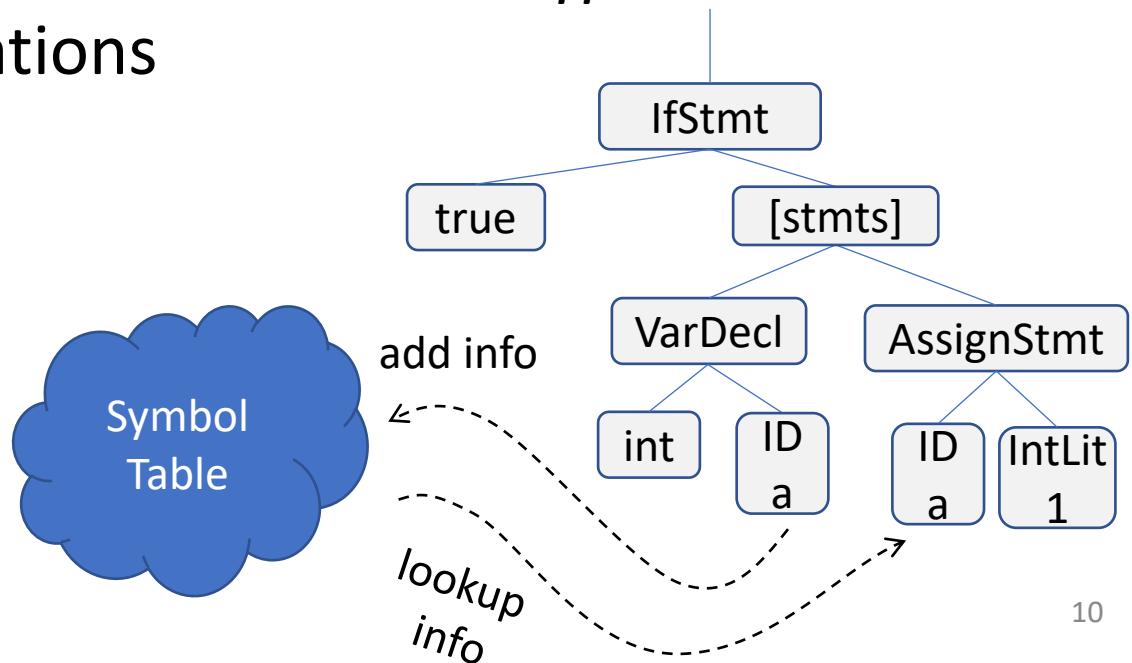
Kinds of entries we need:

- Variable Declarations
- Function Declarations

Program Snippet

```
if(true) {  
    int a;  
    a = 1;  
}
```

AST Snippet



The Compiler's Symbol Table

Semantic Analysis

MyA : A;

What's in the symbol table:

- Depends on the language

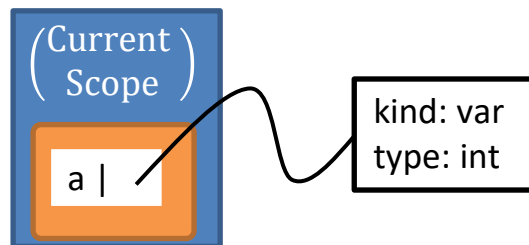
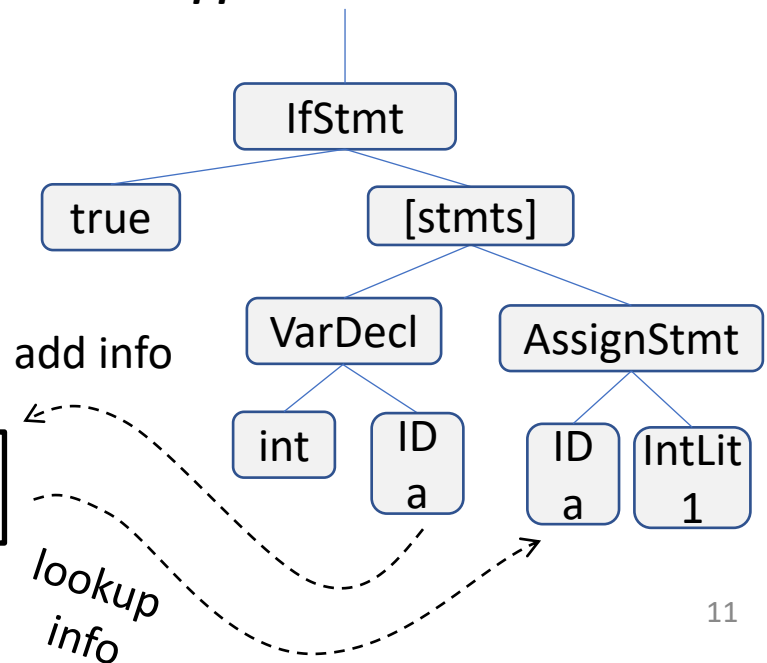
Kinds of entries we need:

- Variable Declarations
- Function Declarations
- Class Declarations

Program Snippet

```
if(true){  
    int a;  
    a = 1;  
}
```

AST Snippet



Symbol Table: A “Snapshot” of Scope

Types – Name Analysis

At any (static) program point

- Symbol table shows *what’s* in scope
- Symbol table shows which scope contains the entry

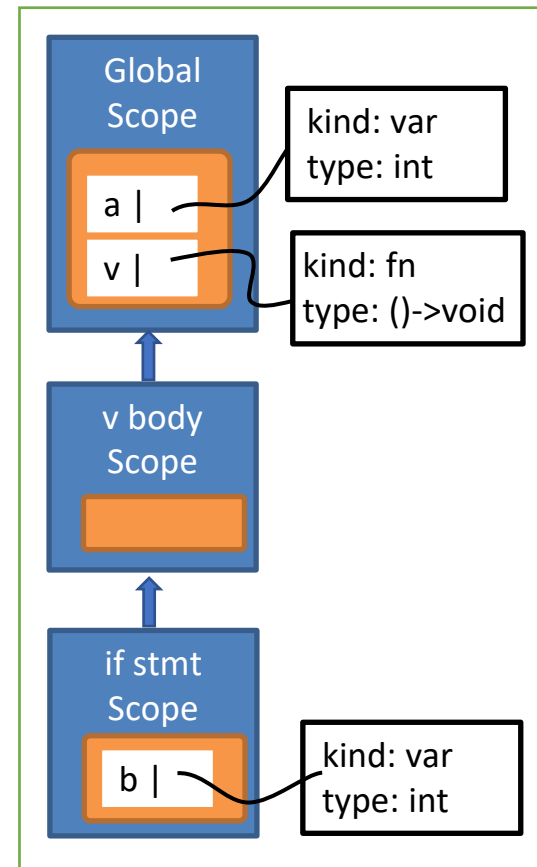
Implementation:

- A list of hashmaps (1 map per scope)

```
1. int a;  
2. void v() {  
3.     if (a) {  
4.         int b;  
5.     }  
6. }
```



Symbol table after line 4

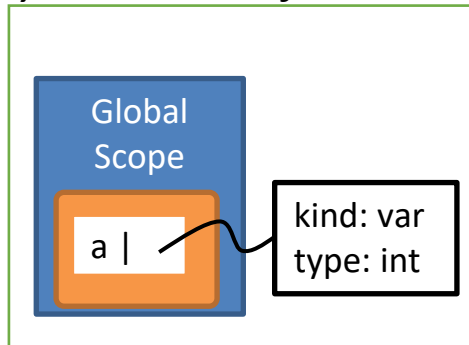


Symbol Table: Scopes “Sub-tables”

Semantic Analysis - Name Analysis

Create one hashmap per scope

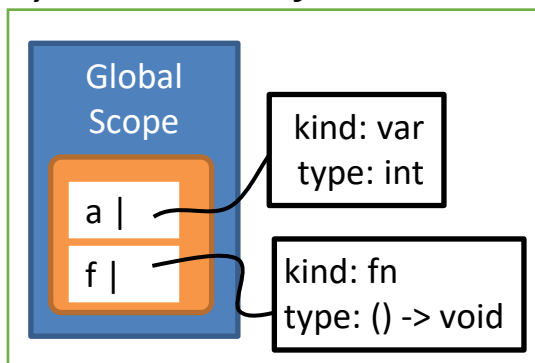
Symbol table after line 1



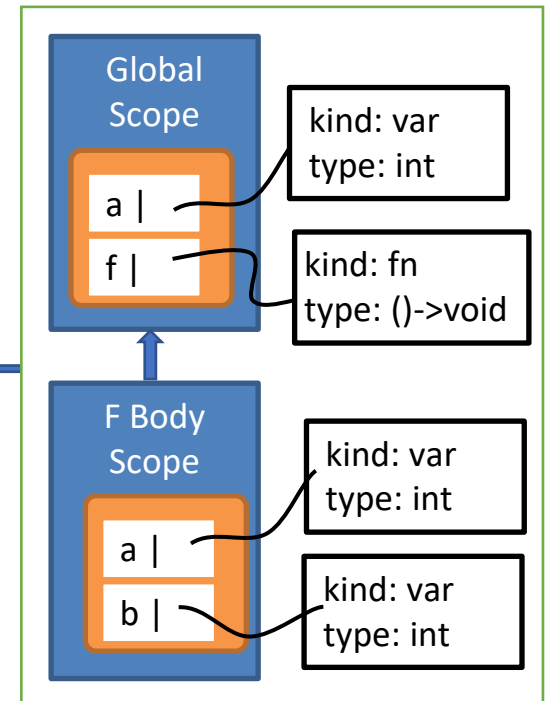
Code

```
1. int a;  
2. void f() {  
3.     int a;  
4.     int b;  
5. }
```

Symbol table after line 5



Symbol table after line 4



Implementation

Semantic Analysis - Name Analysis

- Walk the AST, much like the `unparse()` method
- • Augment AST nodes with a link to the relevant name in the symbol table
- Build new entries into the symbol table when a declaration is encountered
- Connect AST nodes to the entry they add or reference in the symbol table

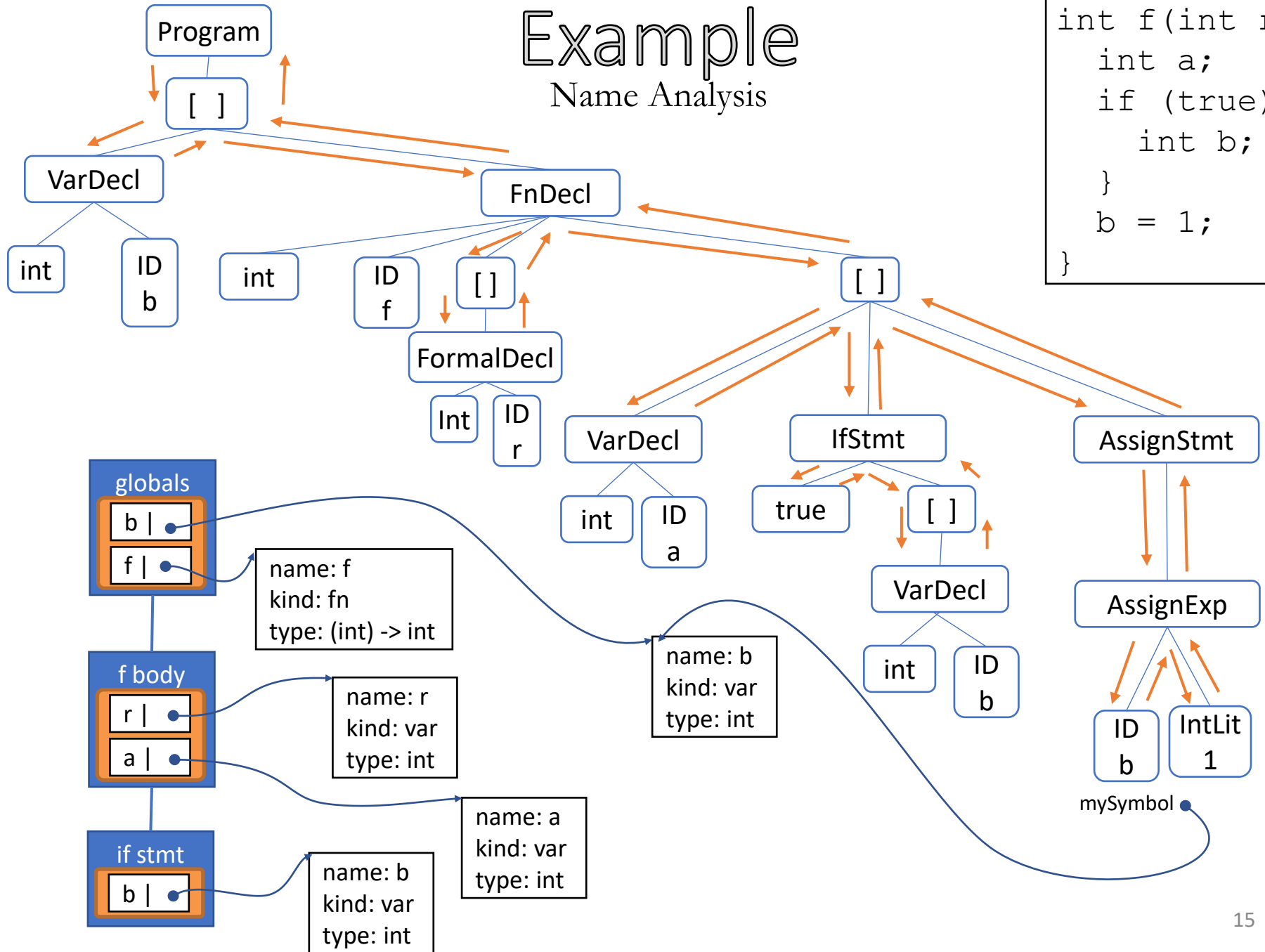
```
class JDNNode {  
    private:  
        String name;  
        Symbol s;  
}
```

new

Example

Name Analysis

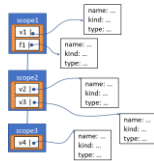
```
int b;
int f(int r){
    int a;
    if (true){
        int b;
    }
    b = 1;
}
```



(My) Terminology

Name Analysis - Implementation

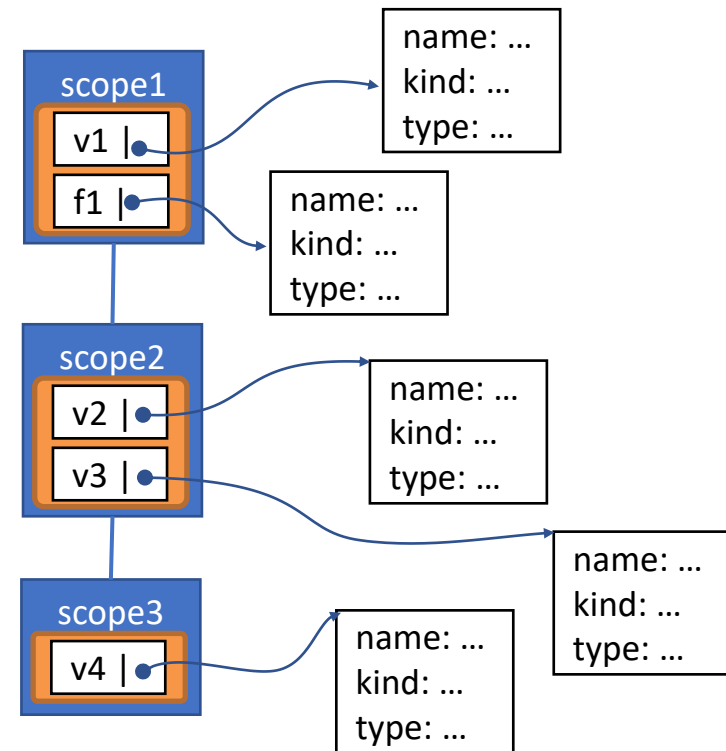
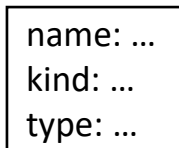
- Symbol Table – the whole structure



- Scope Table – A single map



- Symbol Table Entry (AKA “Semantic Symbol”)



Implementation Suggestions

Name Analysis - Implementation

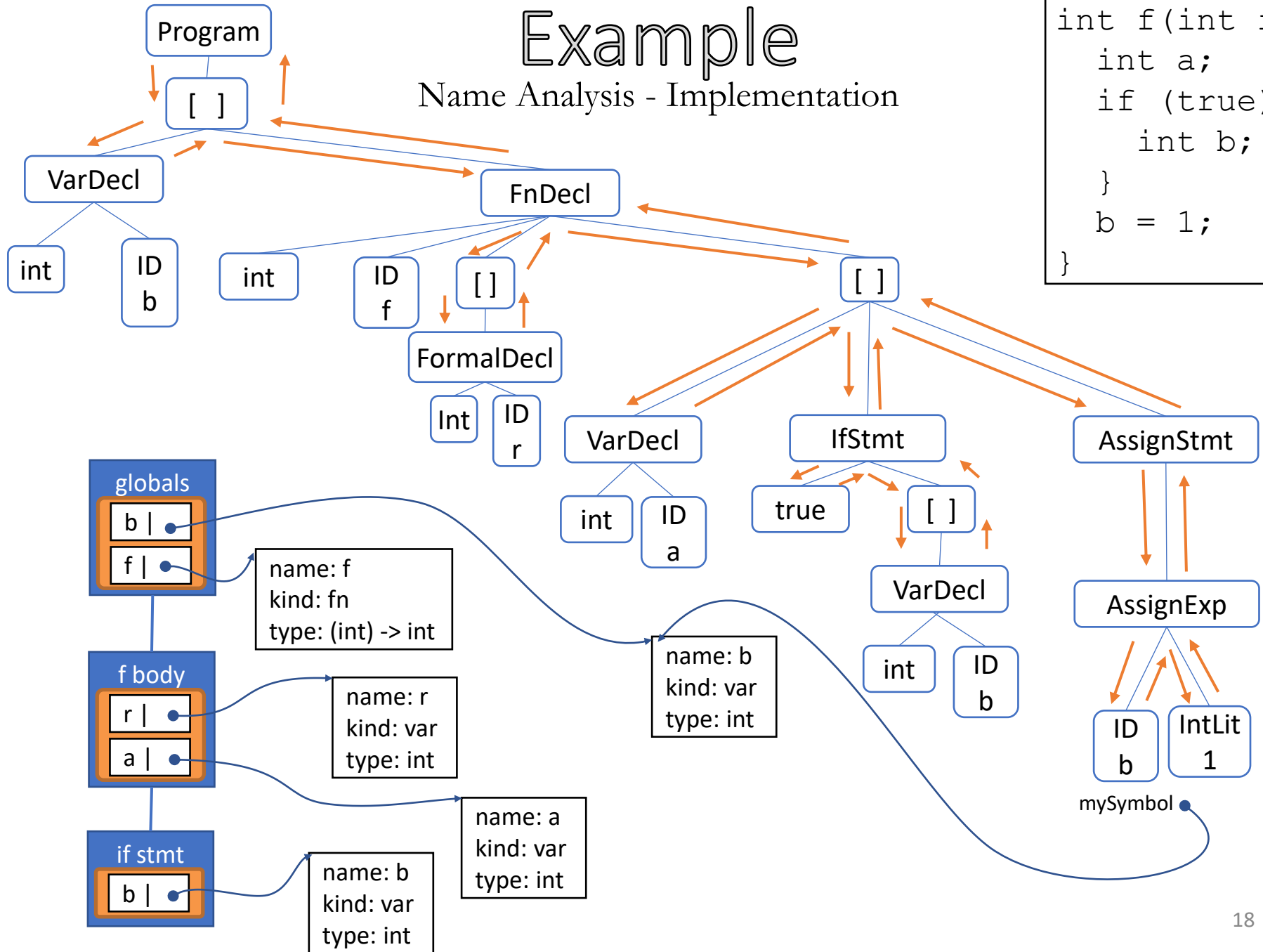
Two approaches

- A nameAnalysis method for each ASTNode subclass
 - Override as appropriate
- The use of the visitor pattern

Example

Name Analysis - Implementation

```
int b;
int f(int r){
    int a;
    if (true){
        int b;
    }
    b = 1;
}
```



Summary

Name Analysis – Wrap-Up

- Described an analysis for enforcing static scoping
- Demonstrated a way to implement the analysis as a walk over the AST

Next Time

Name Analysis – Wrap-Up

Type Systems

- What type systems are
- Why we use them
- The type system for our language

Scratch Page

Name Analysis – Scratch

Scratch Page

Name Analysis – Scratch

Scratch Page

Name Analysis – Scratch

Scratch Page

Name Analysis – Scratch

Scratch Page

Name Analysis – Scratch